

# Proving an Execution of an Algorithm Correct? The case of Polynomial Factorisation

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Thanks to IPAM at UCLA for prompting this, and many colleagues, especially at

Dagstuhl seminar 23401, for input

Partially supported by EPSRC grant EP/T015713

6 October 2023

Do I believe the output from my (complicated, optimised, unverified) computer algebra system?

See JHD's paper at CICM 2023 [Dav23], but note that the same question, in different settings, was asked by Mehlhorn [Meh11] in 1999.

[Dav23] looked at three examples.

**Polynomial Factorisation**  $f = f_1 f_2 \cdots f_k$  and the  $f_i$  is irreducible.

**Integration** The assertion “unintegrable” is correct.

**Real Algebraic Geometry** The assertion that the semi-algebraic variety is empty (UNSAT) is correct.

The last is the most important question, but factorisation is easy to explain and a good case study in its own right.

# Polynomial Factorisation

The base case is polynomials in  $\mathbf{Z}[x]$ .

## Problem (Factorisation)

Given  $f \in \mathbf{Z}[x]$ , write  $f = \prod f_i$  where the  $f_i$  are irreducible elements of  $\mathbf{Z}[x]$ .

Verifying that  $f = \prod f_i$  is, at least relatively, easy. The hard part is verifying that the  $f_i$  are *irreducible*. JHD knows of no implementation of polynomial factorisation that produces any evidence, let alone a proof, of this.

We may as well assume  $f$  is square-free (this would be a rather separate verification question).

# Algorithm

The basic algorithm goes back to [Zas69]: step M is a later addition [Mus75], and the H' variants are also later.

- 1 Choose a prime  $p$  (not dividing the leading coefficient of  $f$ ) such that  $f \pmod{p}$  is also square-free.
  - 2 Factor  $f$  modulo  $p$  as  $\prod f_i^{(1)} \pmod{p}$ .
- M Take five  $p$  and compare the factorisations.
- 3 If  $f$  can be shown to be irreducible from modulo  $p$  factorisations, return  $f$ .
  - 4 Let  $B$  be such that any factor of  $f$  has coefficients less than  $B$  in magnitude, and  $n$  such that  $p^n \geq 2B$ . [Landau–Mignotte]
  - 5 Use Hensel's Lemma to lift the factorisation to  $f = \prod f_i^{(n)} \pmod{p^n}$
- H Starting with singletons and working up, take subsets of the  $f_i^{(n)}$ , multiply them together and check whether, regarded as polynomials over  $\mathbf{Z}$  with coefficients in  $[-B, B]$ , they divide  $f$  — if they do, declare that they are irreducible factors of  $f$ .

H' Use some alternative technique, originally [LLL82], but now e.g. [ASZ00, HvHN11] to find the true factor corresponding to  $f_1^{(n)}$ , remove  $f_1^{(n)}$  and the other  $f_i^{(n)}$  corresponding to this factor, and repeat.



In practice, there are a lot of optimisations, which would greatly complicate a proof of correctness of an implementation of this algorithm.

*We found that, although the Hensel construction is basically neat and simple in theory, the fully optimised version we finally used was as nasty a piece of code to write and debug as any we have come across [MN81].*

Since if  $f$  is irreducible modulo  $p$ , it is irreducible over the integers, the factors produced from singletons in step 5 are easily proved to be irreducible. Unfortunately, the chance that an irreducible polynomial of degree  $n$  is irreducible modulo  $p$  is  $1/n$ .

# Algorithm Notes

A factorisation algorithm could, even though no known implementation does, relatively easily produce the required information for a proof of irreducibility unless the recombination step is required.

**Note** that *verifying* the Hensel lifting, the “nasty piece” from [MN81] is easy: the factors just have to have the right degrees from the factorisation of  $f \pmod{p}$  and multiply to give  $f \pmod{p^n}$ .



Building test cases for the various edge cases was extremely difficult.

Step [H] is relatively easy to verify: this combination divides and no smaller combination divides. The variants in [H'] are interesting: I have not found an easy route.

If [H'] finds a factor that is a product of  $k$   $p$ -adic factors, then we can use [H] to verify this by checking that the  $2^k - 2$  subsets do not give factors.

But if [H'] says “irreducible”, I know no easy proof.

- 1 We can extract from the implementation in FLINT [tea23] of the algorithm with [H], at essentially no cost, the key data that we believe a verifier would need to confirm the irreducibility.
- 2 But this is not necessarily the most efficient verification.
- 3 We *think* that a more efficient verification would need negligibly more work.
- 4 We haven't built a verification.
- 5 The “hard” theorems are (being) stated (LEAN), but what about the “easy” ones, mappings such as “regarded as polynomials over  $\mathbf{Z}$  with coefficients in  $[-B, B]$ ”?
- 6 Needs more theorem prover input.

But We have discovered improvements to FLINT, and at least one new research question in computer algebra.

+ FLINT also has [H'], but we haven't looked at this yet.

# So what is the certificate?

- 1 The 5(?) Musser primes (or the useful subset)
- 2 The factorisations modulo these
  - \* Need to verify that these are irreducible.
- 3 The chosen  $p$  and  $n$
- 4 The set  $S$  of factors modulo  $p^n$ 
  - \* Need to check they match the mod  $p$  ones, and multiply. We have already proved irreducibility of the mod  $p$  versions
- 5 The partition  $S = \bigcup S_i$  that corresponds to the true factorisation.

Thanks to Tobias Nipkow for asking this explicitly.





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