

# Recent advances in real geometric reasoning

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# History of Quantifier Elimination

- In 1930, Tarski discovered [Tar51] that the (semi-)algebraic theory of  $\mathbf{R}^n$  admitted quantifier elimination

$$\exists x_{k+1} \forall x_{k+2} \dots \Phi(x_1, \dots, x_n) \equiv \Psi(x_1, \dots, x_k)$$

- “Semi” = “allowing  $>$ ,  $\leq$  and  $\neq$  as well as  $=$ ”
- Needed as  $\exists y : x = y^2 \Leftrightarrow x \geq 0$
- The complexity of this was indescribable
- In the sense of not being primitive recursive!
- In 1973, Collins [Col75] discovered a much better way:
- Complexity ( $m$  polynomials, degree  $d$ ,  $n$  variables, coefficient length  $l$ )

$$(2d)^{2^{2n+8}} m^{2^{n+6}} l^3 \quad (1)$$

- Construct a cylindrical algebraic decomposition of  $\mathbf{R}^n$ , sign invariant for every polynomial
- Then read off the answer

# What is a CAD?

A **Cylindrical Algebraic Decomposition (CAD)** is a mathematical object. Defined by Collins who also gave the first algorithm to compute one. A CAD is:

- a **decomposition** meaning a partition of  $\mathbf{R}^n$  into connected subsets called **cells**;
- (semi-) **algebraic** meaning that each cell can be defined by a sequence of polynomial equations and inequations;
- **cylindrical** meaning the cells are arranged in a useful manner — their projections are either equal or disjoint.

In addition, there is (usually) a **sample point** in each cell, and an **index** locating it in the decomposition

# “Read off the answer”

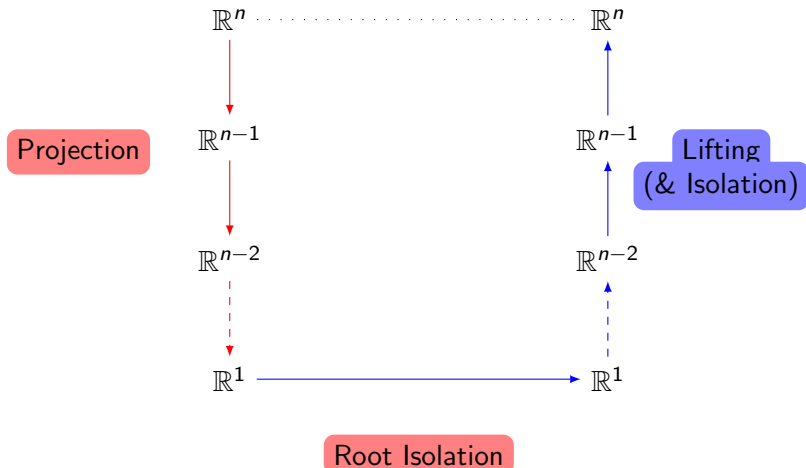
- Each cell is sign invariant, so the the truth of a formula **throughout** the cell is the truth at the sample point.
- $\forall x F(x) \Leftrightarrow$  “ $F(x)$  is true at all sample points”
- $\exists x F(x) \Leftrightarrow$  “ $F(x)$  is true at some sample point”
- $\forall x \exists y F(x, y) \Leftrightarrow$  “take a CAD of  $\mathbf{R}^2$ , cylindrical for  $y$  projected onto  $x$ -space, then check

$\forall$  sample  $x \exists$  sample  $(x, y) : F(x, y)$  is true”: finite check

**NB** The order of the quantifiers defines the order of projection

So all we need is a CAD!

# The basic idea for CAD [Col75]



# So how do we project?

(Lifting is in fact relatively straight-forward)

Given polynomials  $\mathcal{P}_n = \{p_i\}$  in  $x_1, \dots, x_n$ , what should  $\mathcal{P}_{n-1}$  be?

Naïve (Doesn't work!) Every  $\text{disc}_{x_n}(p_i)$ , every  $\text{res}_{x_n}(p_i, p_j)$

i.e. where the polynomials fold, or cross: misses lots of "special" cases

[Col75] First enlarge  $\mathcal{P}_n$  with all its reducta, then naïve plus the coefficients of  $\mathcal{P}_n$  (with respect to  $x_n$ ) the principal subresultant coefficients from the  $\text{disc}_{x_n}$  and  $\text{res}_{x_n}$  calculations

[Hon90] a tidied version of [Col75].

[McC88] Let  $\mathcal{B}_n$  be a squarefree basis for the primitive parts of  $\mathcal{P}_n$ . Then  $\mathcal{P}_{n-1}$  is the contents of  $\mathcal{P}_n$ , the coefficients of  $\mathcal{B}_n$  and every  $\text{disc}_{x_n}(b_i)$ ,  $\text{res}_{x_n}(b_i, b_j)$  from  $\mathcal{B}_n$

[Bro01] Naïve plus leading coefficients (not squarefree!)

# Are these projections correct?

[Col75] Yes, and it's relatively straightforward to prove that, over a cell in  $\mathbf{R}^{n-1}$  sign-invariant for  $\mathcal{P}_{n-1}$ , the polynomials of  $\mathcal{P}_n$  do not cross, and define cells sign-invariant for the polynomials of  $\mathcal{P}_n$

[McC88] 52 pages (based on [Zar75]) prove the equivalent statement, but for **order-invariance**, not sign-invariance, provided the polynomials are **well-oriented**, a test that has to be applied during lifting.

But if they're not known to be well-oriented?

[McC88] suggests adding all partial derivatives

In practice hope for well-oriented, and if it fails use Hong's projection.

[Bro01] Needs well-orientedness and additional checks

# What about the complexity?

If the McCallum projection is well-oriented, the complexity is

$$(2d)^{n2^{n+7}} m^{2^{n+4}} l^3 \quad (2)$$

versus the original

$$(2d)^{2^{2n+8}} m^{2^{n+6}} l^3 \quad (1)$$

and in practice the gains in running time can be factors of a thousand, or, more often, the difference between feasibility and infeasibility

“Randomly”, well-orientedness ought to occur with probability 1, but we have a family of “real-world” examples where it often fails



# Need it be this hard?

The Heintz construction

$$\Phi_k(x_k, y_k) := \left[ \begin{array}{l} \exists z_k \forall x_{k-1} y_{k-1} \left[ \begin{array}{l} y_{k-1} = y_k \wedge x_{k-1} = z_k \vee y_{k-1} = z_k \wedge x_{k-1} = x_k \\ \Rightarrow \Phi_{k-1}(x_{k-1}, y_{k-1}) \end{array} \right] \end{array} \right]$$

If  $\Phi_1 \equiv y_1 = f(x_1)$ , then  $\Phi_2 \equiv y_2 = f(f(x_2))$ ,

$\Phi_3 \equiv y_3 = f(f(f(f(x_3))))$

[DH88] shows  $\Omega\left(2^{2^{(n-2)/5}}\right)$  (using  $y_R + iy_I = (x_R + ix_I)^4$ )

[BD07] shows  $\Omega\left(2^{2^{(n-1)/3}}\right)$  (using a sawtooth)

Hence doubly exponential is inevitable, but there's a lot of room!

In fact, there are theoretical algorithms which are singly-exponential in  $n$ , but doubly-exponential in the number of  $\exists \forall$  alternations

[McC99] “equational constraints” : when

$$\Phi \equiv f(x, y, \dots) = 0 \wedge (\dots)$$

Note If  $\Phi \equiv (f_1(x, y) = 0 \wedge g_1(x, y) < 0) \vee (f_2(x, y) = 0 \wedge g_2(x, y) < 0)$ , which has no obvious equational constraint, we can consider  $(f_1 \cdot f_2)(x, y) = 0 \wedge \Phi$ , which is equivalent (but higher degree)

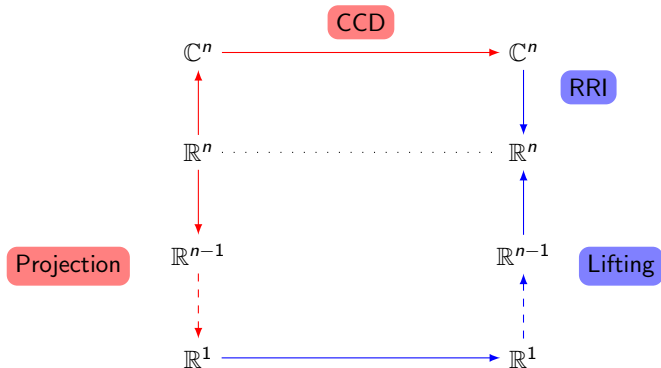
[BDE<sup>+</sup>13] “truth table invariant CAD” treats this directly

submitted also handles the case where not every clause has an equality

Roughly speaking, the effect is to reduce  $n$  by 1, which square roots the complexity

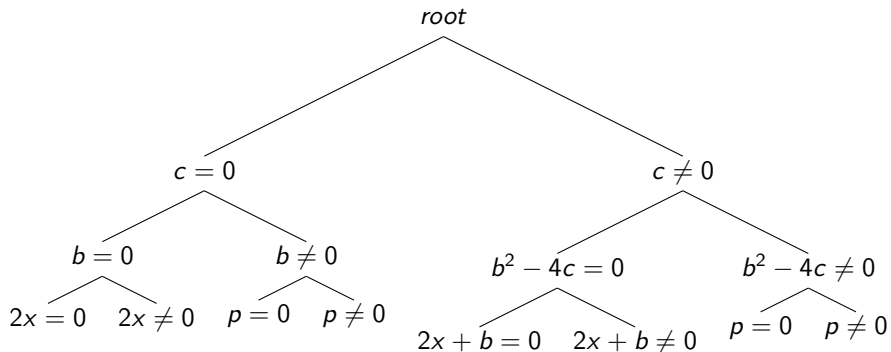
# An alternative approach [CMMXY09]

Proceed via the complex numbers,



Do a complex cylindrical decomposition via **Regular Chains**  
Can be combined with truth table ideas [BCD<sup>+</sup>14]

# Example Complex CD



**Figure:** Complete complex cylindrical tree for the general monic quadratic equation,  $p := x^2 + bx + c$ , under variable ordering  $c \prec b \prec x$ .

Note that  $b = 0$  is only tested where relevant

# So might I trust these results?

Trivially for  $\exists$  problems a positive result, or negative for  $\forall$  problems, is easily verified (witness computation)

Negative  $\exists$  is essentially refutation [JdM12]

Otherwise we're believing a complicated software package and some maths

[Col75] Algebra system + 3200LOC + “some maths”

[McC88] Algebra system + 3200LOC + “a lot of maths”

[CMMXY09] Algebra system + 5000LOC + “medium maths”

[BDE<sup>+</sup>13] Algebra system + 6200LOC + “medium maths”

Proven software? [CM12] does QE (not full CAD), loosely based on [Col75], in COQ, but terribly impractical

Note that CAD has other applications — algebraic simplification [BCD<sup>+</sup>02], robot path planning [SS83], which tends to require adjacency(unsolved in general dimension)



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