# Preconditioned inexact inverse iteration and inexact shift-invert Arnoldi method 

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Preparation for ITT2
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1 Introduction

2 Inexact inverse iteration

3 Inexact Shift-invert Arnoldi method

4 Conclusions

## Outline

1 Introduction

2 Inexact inverse iteration

3 Inexact Shift-invert Arnoldi method

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## Motivation

Alastairs lecture:

$$
\left(I-\sigma_{s} K_{\sigma}\right) \Phi=\lambda \nu \sigma_{f} K_{\sigma} \Phi
$$

■ Eigenvalue problem arising from reactor criticality computation

- Theory on the derivation in the thesis by Fynn Scheben
- Interest in the smallest eigenvalue (which is real and positive)


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■ Can be written as

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A x=\lambda B x, \quad \lambda \in \mathbb{C}, x \in \mathbb{C}^{n}
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$$

- In this talk: $B=I$ (but all results extend to $B \neq I$ )


## Problem and iterative methods

Find a small number of eigenvalues and eigenvectors of:

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■ $A$ is large, sparse, nonsymmetric

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- Power method
- Simultaneous iteration
- Arnoldi method
- Jacobi-Davidson method
- repeated application of the matrix $A$ to a vector

■ Generally convergence to largest/outlying eigenvector

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## Shift-invert strategy

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## Shift-invert strategy

- Wish to find a few eigenvalues close to a shift $\sigma$

- Problem becomes

$$
(A-\sigma I)^{-1} x=\frac{1}{\lambda-\sigma} x
$$

$\square$ each step of the iterative method involves repeated application of $\mathcal{A}=(A-\sigma I)^{-1}$ to a vector

- Inner iterative solve:

$$
(A-\sigma I) y=x
$$

using Krylov method for linear systems. (CG, MINRES, GMRES, ...)
■ leading to inner-outer iterative method.

## Shift-invert strategy

# This talk: <br> Convergence of outer iteration 

Inner iteration and preconditioning

Inverse iteration and Arnoldi method

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## Classical methods for finding one eigenvalue

## Inverse iteration/Rayleigh quotient iteration

INPUT: Matrix $A$ shift $\sigma \approx \lambda$, initial vector $x_{0}$ with $\left\|x_{0}\right\|=1$. OUTPUT: Approximate eigenpair $(\lambda, x)$ for $i=1$ to $\ldots$ do

Choose shift $\sigma$.
Solve for $\hat{x}$

$$
(A-\sigma I) \hat{x}=x
$$

Rescale $x=\frac{\hat{x}}{\|\hat{x}\|}$,
Update $\lambda=\rho(x)=x^{H} A x$,
Test for convergence (using eigenvalue residual $(r=(A-\lambda I) x)$.
end for

## Classical methods for finding one eigenvalue

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OUTPUT: Approximate eigenpair $(\lambda, x)$ for $i=1$ to $\ldots$ do

Choose shift $\sigma$.
Run $k$ steps of a Krylov subspace method to obtain $\hat{x}$ such that

$$
0 \leq\|(A-\sigma I) \hat{x}-x\| \leq \xi
$$

Rescale $x=\frac{\hat{x}}{\|\hat{x}\|}$,
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## Classical methods for finding one eigenvalue

## Convergence rates of exact methods

■ Inverse iteration: linear convergence
■ Rayleigh quotient iteration (RQI): cubic convergence for normal $A$ (otherwise quadratic)

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## Convergence rate analysis of inexact methods

Main requirement: decreasing accuracy of the inner solves:

$$
\xi^{(i+1)} \leq \xi^{(i)}
$$

The convergence speed of the exact methods can be re-established.
[Lai/Lin/Lin '97, Golub/Ye '00, Simoncini/Elden '02, F./Spence '07, Elman/Xue '09, ...]

## Error indicator

## Error indicator (Orthogonal decomposition, symmetric A)


$\mathrm{Qx}{ }^{(i)}=\mathrm{O}\left(\sin \theta^{(i)}\right)$ measure for the error

$$
x^{(i)}=\cos \theta^{(i)} x_{1}+\sin \theta^{(i)} x_{\perp}^{(i)}, \quad x_{\perp}^{(i)} \perp x_{1} .
$$

Generalisations to nonsymmetric $A$ exist.

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Generalisations to nonsymmetric $A$ exist.

Eigenvalue residual

$$
C_{1}\left|\sin \theta^{(i)}\right| \leq\left\|r^{(i)}\right\| \leq C_{2}\left|\sin \theta^{(i)}\right|
$$ solver

Error in $\theta$ (eigenvector)

$$
\tan \theta^{(i+1)} \leq \frac{\left|\lambda_{1}-\sigma^{(i)}\right|}{\left|\lambda_{2}-\sigma^{(i)}\right|} T^{(i)}
$$

■ Exact solves: $T^{(i)}=\tan \theta^{(i)}$
$\square$ Inexact solves $T^{(i)}=\frac{\left|\sin \theta^{(i)}\right|+\left\|\left(I-x_{1} x_{1}^{*}\right) s^{(i)}\right\|}{\mid \cos \theta^{(i)}-\left\|x_{1}^{*} s^{(i)}\right\|}, \quad\left\|s^{(i)}\right\| \leq \xi^{(i)}$ solver

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## Choice of $\tau$

■ Choice 1: $\xi^{(i)}=C\left\|r^{(i)}\right\|=\mathcal{O}\left(\sin \theta^{(i)}\right) \Rightarrow T^{(i)}=\mathcal{O}\left(\tan \theta^{(i)}\right)$
■ Choice 2: $\xi^{(i)}=\mathrm{constant} \Rightarrow T^{(i)}=\mathcal{O}(1)$

## Classical methods for finding one eigenvalue

## Inexact Inverse iteration/Inexact Rayleigh quotient iteration

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OUTPUT: Approximate eigenpair $(\lambda, x)$
for $i=1$ to $\ldots$ do
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Run $k$ steps of a Krylov subspace method to obtain $\hat{x}$ such that

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Rescale $x=\frac{\hat{x}}{\|\hat{x}\|}$,
Update $\lambda=\rho(x)=x^{H} A x$,
Test for convergence (using eigenvalue residual $(r=(A-\lambda I) x)$. end for

## Convergence rates

If the solve tolerance is decreased, i.e.

$$
\xi^{(i)}=C\left\|r^{(i)}\right\|
$$

then convergence rate is linear (same convergence rate as for exact solves).

## Inexact methods

Example: matrix sherman5, $n=3312$, fixed shift $\sigma=0$.

- exact solution (Matlab backslash)
- inexact solution (full GMRES) with decreasing tolerance
- termination when $\left\|r^{(i)}\right\| \leq 10^{-9}$, where
- preconditioned GMRES with incomplete LU preconditioners $P$


## Inexact methods

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## Inexact methods - dilemma!

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increasing accuracy $\Rightarrow$ number of inner iterations remains $\approx$ constant (fixed accuracy $\Rightarrow$ number of inner iterations $\approx$ decreasing)

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$\square$ right hand side $x$ is an approximate eigenvector in inexact RQI
■ beneficial for Krylov subspace methods!
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## Basic idea - key observation

Solving $B x=b$ with a Krylov method, where $b$ is an eigenvector of $B$ :

$$
\begin{aligned}
\Rightarrow \mathcal{K}_{1}(B, b) & =\operatorname{span}\{b\} \\
\mathcal{K}_{2}(B, b) & =\operatorname{span}\{b, B b\}=\operatorname{span}\{b, \alpha b\}=\operatorname{span}\{b\}
\end{aligned}
$$

Krylov method converges in one step!

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preconditioning $\Rightarrow$ number of inner iterations is clearly increasing

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■ Solution of $(A-\sigma I) \hat{x}=x$ with a (left) preconditioned Krylov subspace method:

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P^{-1}(A-\sigma I) \hat{x}=P^{-1} x
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but $P^{-1} x$ is a poor eigenvector approximation of $P^{-1}(A-\sigma I)$ !.

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- This is also the case for the generalised eigenproblem where we solve

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$$

## Question

Can we reduce the number of inner steps, e.g. to an $\approx$ constant level?

## Tuned preconditioners

We need a preconditioner $\mathbb{P}$ such that $\mathbb{P}^{-1} x$ is an approximate eigenvector of $\mathbb{P}^{-1}(A-\sigma I)$ :

$$
\mathbb{P}^{-1}(A-\sigma I) \mathbb{P}^{-1} x=(\lambda-\sigma) \mathbb{P}^{-1} x
$$

Has to satisfy $\mathbb{P} x=x($ or $\mathbb{P} x=A x)$

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Has to satisfy $\mathbb{P} x=x($ or $\mathbb{P} x=A x)$

## The tuned preconditioner for one-sided RQI

For $P \approx A$, the tuned preconditioner $\mathbb{P}$ is defined by

$$
\mathbb{P}=P+(x-P x) x^{H}
$$

and we obtain

$$
\mathbb{P}^{-1}=P^{-1}-\frac{\left(P^{-1} x-x\right) x^{H} P^{-1}}{x^{H} P^{-1} x}
$$

Minor modification and minor extra computational cost.

## Inexact two-sided methods - dilemma!

Example: matrix sherman5, $n=3312$, fixed shift $\sigma=0$.


## Inexact two-sided methods - dilemma resolved!

Example: matrix sherman5, $n=3312$, fixed shift $\sigma=0$.


## Standard GMRES theory for $y_{0}=0$ and $A$ diagonalisable

$$
\left\|x-(A-\sigma I) y_{k}\right\| \leq \kappa(W) \min _{p \in \mathcal{P}_{k}} \max _{j=1, \ldots, n} \mid p\left(\lambda_{j}\right)\|x\|
$$

where $\lambda_{j}$ are eigenvalues of $A-\sigma I$ and $(A-\sigma I)=W \Lambda W^{-1}$.

The inner iteration for $(A-\sigma I) y=x$

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$$

where $\lambda_{j}$ are eigenvalues of $A-\sigma I$ and $(A-\sigma I)=W \Lambda W^{-1}$.
Number of inner iterations

$$
k \geq C_{1}+C_{2} \log \frac{\|x\|}{\xi}
$$

for $\left\|x-(A-\sigma I) y_{k}\right\| \leq \xi$.

## More detailed GMRES theory for $y_{0}=0$

$$
\left\|x-(A-\sigma I) y_{k}\right\| \leq \tilde{\kappa}(W) \frac{\left|\lambda_{2}-\lambda_{1}\right|}{\lambda_{1}} \min _{p \in \mathcal{P}_{k-1}} \max _{j=2, \ldots, n}\left|p\left(\lambda_{j}\right)\right|\|\mathcal{Q} x\|
$$

where $\lambda_{j}$ are eigenvalues of $A-\sigma I$.

## Number of inner iterations

$$
k \geq C_{1}^{\prime}+C_{2}^{\prime} \log \frac{\|\mathcal{Q} x\|}{\xi}
$$

where $\mathcal{Q}$ projects onto the space not spanned by the eigenvector.

## The inner iteration for $(A-\sigma I) y=x$

## Good news!

$$
k^{(i)} \geq C_{1}^{\prime}+C_{2}^{\prime} \log \frac{C_{3}\left\|r^{(i)}\right\|}{\xi^{(i)}}
$$

where $\xi^{(i)}=C\left\|r^{(i)}\right\|$. Iteration number approximately constant!

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where $\xi^{(i)}=C\left\|r^{(i)}\right\|$. Iteration number approximately constant!

## Bad news :-(

For a standard preconditioner $P$

$$
\begin{gathered}
(A-\sigma I) P^{-1} \tilde{y}^{(i)}=x^{(i)} \quad P^{-1} \tilde{y}^{(i)}=y^{(i)} \\
k^{(i)} \geq C_{1}^{\prime \prime}+C_{2}^{\prime \prime} \log \frac{\left\|\tilde{\mathcal{Q}} x^{(i)}\right\|}{\xi^{(i)}}=C_{1}^{\prime \prime}+C_{2}^{\prime \prime} \log \frac{C}{\xi^{(i)}}
\end{gathered}
$$

where $\xi^{(i)}=C\left\|r^{(i)}\right\|$. Iteration number increases!

## Convection-Diffusion operator

Finite difference discretisation on a $32 \times 32$ grid of the convection-diffusion operator

$$
-\Delta u+5 u_{x}+5 u_{y}=\lambda u \quad \text { on } \quad(0,1)^{2}
$$

with homogeneous Dirichlet boundary conditions ( $961 \times 961$ matrix).
■ smallest eigenvalue: $\lambda_{1} \approx 32.18560954$,
■ Preconditioned GMRES with tolerance $\xi^{(i)}=0.01\left\|r^{(i)}\right\|$,
■ standard and tuned preconditioner (incomplete LU).

## Convection-Diffusion operator



Figure: Inner iterations vs outer iterations


Figure: Eigenvalue residual norms vs total number of inner iterations

## Convection-Diffusion operator



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## The algorithm

## Arnoldi's method

- Arnoldi method constructs an orthogonal basis of $k$-dimensional Krylov subspace

$$
\begin{gathered}
\mathcal{K}_{k}\left(\mathcal{A}, q^{(1)}\right)=\operatorname{span}\left\{q^{(1)}, \mathcal{A} q^{(1)}, \mathcal{A}^{2} q^{(1)}, \ldots, \mathcal{A}^{k-1} q^{(1)}\right\}, \\
\mathcal{A} Q_{k}=Q_{k} H_{k}+q_{k+1} h_{k+1, k} e_{k}^{H}=Q_{k+1}\left[\begin{array}{c}
H_{k} \\
h_{k+1, k} e_{k}^{H}
\end{array}\right] \\
Q_{k}^{H} Q_{k}=I .
\end{gathered}
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■ Eigenvalues of $H_{k}$ are eigenvalue approximations of (outlying) eigenvalues of $\mathcal{A}$

$$
\left\|r_{k}\right\|=\|\mathcal{A} x-\theta x\|=\left\|\left(\mathcal{A} Q_{k}-Q_{k} H_{k}\right) u\right\|=\left|h_{k+1, k} \| e_{k}^{H} u\right|
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$$

■ at each step, application of $\mathcal{A}$ to $q_{k}: \mathcal{A} q_{k}=\tilde{q}_{k+1}$

## Example

random complex matrix of dimension $n=144$ generated in Matlab: $G=$ numgrid(' $N$ ', 14) ; $B=\operatorname{delsq}(G) ; A=\operatorname{sprandn}(B)+i * \operatorname{sprandn}(B)$


## after 5 Arnoldi steps



## after 10 Arnoldi steps



## after 15 Arnoldi steps



## after 20 Arnoldi steps



## after 25 Arnoldi steps



## after 30 Arnoldi steps



## Shift-Invert Arnoldi's method $\mathcal{A}:=A^{-1}$

- Arnoldi method constructs an orthogonal basis of $k$-dimensional Krylov subspace

$$
\begin{gathered}
\mathcal{K}_{k}\left(A^{-1}, q^{(1)}\right)=\operatorname{span}\left\{q^{(1)}, A^{-1} q^{(1)},\left(A^{-1}\right)^{2} q^{(1)}, \ldots,\left(A^{-1}\right)^{k-1} q^{(1)}\right\}, \\
A^{-1} Q_{k}=Q_{k} H_{k}+q_{k+1} h_{k+1, k} e_{k}^{H}=Q_{k+1}\left[\begin{array}{c}
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- at each step, application of $A^{-1}$ to $q_{k}: A^{-1} q_{k}=\tilde{q}_{k+1}$

Inexact solves (Simoncini 2005), Bouras and Frayssé (2000)

- Wish to solve

$$
\left\|q_{k}-A \tilde{q}_{k+1}\right\|=\left\|\tilde{d}_{k}\right\| \leq \tau_{k}
$$

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\end{array}\right]+\left[d_{1}|\ldots| d_{k}\right]
$$

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- $u$ eigenvector of $H_{k}$ :

$$
\left\|r_{k}\right\|=\left\|\left(A^{-1} Q_{k}-Q_{k} H_{k}\right) u\right\|=\left|h_{k+1, k} \| e_{k}^{H} u\right|+D_{k} u
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$$

- Linear combination of the columns of $D_{k}$

$$
D_{k} u=\sum_{l=1}^{k} d_{l} u_{l}, \quad \text { if } \quad u_{l} \quad \text { small, then } \quad\left\|d_{l}\right\| \quad \text { allowed to be large! }
$$

## Inexact solves (Simoncini 2005), Bouras and Frayssé (2000)

Linear combination of the columns of $D_{k}$

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\begin{gathered}
D_{k} u=\sum_{l=1}^{k} d_{l} u_{l}, \quad \text { if } u_{l} \quad \text { small, then }\left\|d_{l}\right\| \quad \text { allowed to be large! } \\
\left\|d_{l} u_{l}\right\| \leq \frac{1}{k} \varepsilon \Rightarrow\left\|D_{k} u\right\|<\varepsilon \\
\left|u_{l}\right| \leq C(l, k)\left\|r_{l-1}\right\|
\end{gathered}
$$

and
leads to

$$
\begin{aligned}
& \left\|q_{k}-A \tilde{q}_{k+1}\right\|=\left\|\tilde{d}_{k}\right\| \\
& \left\|\tilde{d}_{k}\right\|=C \frac{1}{\left\|r_{k-1}\right\|}
\end{aligned}
$$

Solve tolerance can be relaxed.

## Preconditioning

GMRES convergence bound

$$
\left\|q_{k}-A P^{-1} \tilde{q}_{k+1}^{l}\right\|=\kappa \min _{p \in \Pi_{l}} \max _{i=1, \ldots, n}\left|p\left(\mu_{i}\right)\right|\left\|q_{k}\right\|
$$

depending on

## The inner iteration for $A P^{-1} \tilde{q}_{k+1}=q_{k}$

## Preconditioning

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$$

depending on

- the eigenvalue clustering of $A P^{-1}$
- the condition number
- the right hand side (initial guess)


## Preconditioning

- Introduce preconditioner $P$ and solve

$$
A P^{-1} \tilde{q}_{k+1}=q_{k}, \quad P^{-1} \tilde{q}_{k+1}=q_{k+1}
$$

using GMRES

## Preconditioning

- Introduce preconditioner $P$ and solve

$$
A P^{-1} \tilde{q}_{k+1}=q_{k}, \quad P^{-1} \tilde{q}_{k+1}=q_{k+1}
$$

using GMRES

## Tuned Preconditioner

using a tuned preconditioner for Arnoldi's method

$$
\mathbb{P}_{k} Q_{k}=A Q_{k} ; \quad \text { given by } \quad \mathbb{P}_{k}=P+(A-P) Q_{k} Q_{k}^{H}
$$

## Theorem (Properties of the tuned preconditioner)

Let $P$ with $P=A+E$ be a preconditioner for $A$ and assume $k$ steps of Arnoldi's method have been carried out; then $k$ eigenvalues of $A \mathbb{P}_{k}^{-1}$ are equal to one:

$$
\left[A \mathbb{P}_{k}^{-1}\right] A Q_{k}=A Q_{k}
$$

and $n-k$ eigenvalues are close to the corresponding eigenvalues of $A P^{-1}$.

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## Implementation

■ Sherman-Morrison-Woodbury.
■ Only minor extra costs (one back substitution per outer iteration)

## Numerical Example

sherman5.mtx nonsymmetric matrix from the Matrix Market library $(3312 \times 3312)$.

■ smallest eigenvalue: $\lambda_{1} \approx 4.69 \times 10^{-2}$,

- Preconditioned GMRES as inner solver (both fixed tolerance and relaxation strategy),
■ standard and tuned preconditioner (incomplete LU).


## No tuning and standard preconditioner



Figure: Inner iterations vs outer iterations


Figure: Eigenvalue residual norms vs total number of inner iterations

## Tuning the preconditioner



Figure: Inner iterations vs outer iterations


Figure: Eigenvalue residual norms vs total number of inner iterations

## Relaxation



Figure: Inner iterations vs outer iterations

Figure: Eigenvalue residual norms vs total number of inner iterations


Figure: Inner iterations vs outer iterations


Figure: Eigenvalue residual norms vs total number of inner iterations

## Ritz values of exact and inexact Arnoldi

| Exact eigenvalues | Ritz values (exact Arnoldi) | Ritz values (inexact Arnoldi, tuning) |
| :---: | :---: | :---: |
| $+4.69249563 \mathrm{e}-02$ | $+\underline{4.69249563 \mathrm{e}-02}$ | $+\underline{4.69249563 \mathrm{e}-02}$ |
| $+1.25445378 \mathrm{e}-01$ | $+\underline{1.25445378} \mathrm{e}-01$ | $+\underline{1.25445378} \mathrm{e}-01$ |
| $+4.02658363 \mathrm{e}-01$ | $+\underline{4.02658347} \mathrm{e}-01$ | $+\underline{4.02658244} \mathrm{e}-01$ |
| $+5.79574381 \mathrm{e}-01$ | $+\underline{5.79} 625498 \mathrm{e}-01$ | $+\underline{5.79817301 \mathrm{e}-01}$ |
| $+6.18836405 \mathrm{e}-01$ | $+\underline{6.18} 798666 \mathrm{e}-01$ | $+\underline{6.18} 650849 \mathrm{e}-01$ |

Table: Ritz values of exact Arnoldi's method and inexact Arnoldi's method with the tuning strategy compared to exact eigenvalues closest to zero after 14 shift-invert Arnoldi steps.

## Outline

1 Introduction

2 Inexact inverse iteration

3 Inexact Shift-invert Arnoldi method

4 Conclusions

## Conclusions

■ For eigenvalue computations it is advantageous to consider small rank changes to the standard preconditioners

- Works for any preconditioner
- Works for SI versions of Power method, Simultaneous iteration, Arnoldi method
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