

# Dynamics in geometric scale-free networks

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based on joint work with *Arne Grauer*

**Stochastic Geometry in Action**

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# Structure of talk

The weight-dependent random connection model

The contact process on the weight-dependent random connection model

The dynamic weight-dependent random connection model

## **The weight-dependent random connection model**

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- **Vertices:** a Poisson point process  $\mathcal{X}$  of unit intensity on  $\mathbb{R}^d \times [0, 1]$  or  $\mathbb{T}_n^d \times [0, 1]$ .

# The weight-dependent random connection model

- **Vertices:** a Poisson point process  $\mathcal{X}$  of unit intensity on  $\mathbb{R}^d \times [0, 1]$  or  $\mathbb{T}_n^d \times [0, 1]$ .
- **Edges:** Given  $\mathcal{X}$ , edges are drawn independently of each other and there exists  $\alpha, \kappa_1, \kappa_2 > 0$  such that, for every pair of vertices  $\mathbf{x} = (x, t), \mathbf{y} = (y, s) \in \mathcal{X}$ , it holds

## Assumption (A1)

$$\begin{aligned} & \alpha (1 \wedge \kappa_1 (t \wedge s)^{-\delta\gamma} |x - y|^{-\delta d}) \\ & \leq \mathbb{P}_{\mathbf{x}, \mathbf{y}}(\mathbf{x} \sim \mathbf{y}) \leq \\ & \kappa_2 (t \wedge s)^{-\delta\gamma} (t \vee s)^{\delta(\gamma-1)} |x - y|^{-\delta d}, \end{aligned}$$

where  $\gamma \in (0, 1)$ ,  $\delta > 1$ ,  $\alpha, \kappa_1, \kappa_2 > 0$ .

## Examples of the weight-dependent random connection model

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Let  $\varphi : (0, \infty) \rightarrow [0, 1]$  be a non-decreasing and satisfying  $\varphi(r) \asymp r^{-\delta}$ . Form edges between  $(x, t)$  and  $(y, s)$  with probability

$$\varphi(\beta^{-1}(t \wedge s)^\gamma(t \vee s)^{1-\gamma}|x - y|^d).$$

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- **The soft Boolean model:**

Let the radius distribution satisfy  $\mathbb{P}(R_x > r) \asymp r^{-d/\gamma}$  as  $r \rightarrow \infty$  and let  $X(x, y)$  satisfy

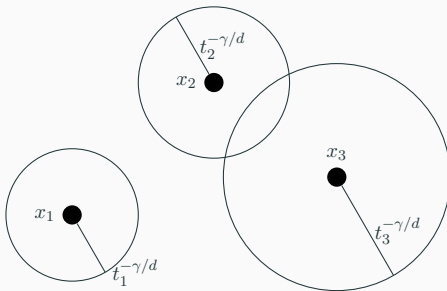
$$\mathbb{P}(X > r) \asymp r^{-\delta d} \text{ as } r \rightarrow \infty.$$

Connect to vertices when  $\frac{|x-y|}{R_x+R_y} \leq X(x, y)$ .



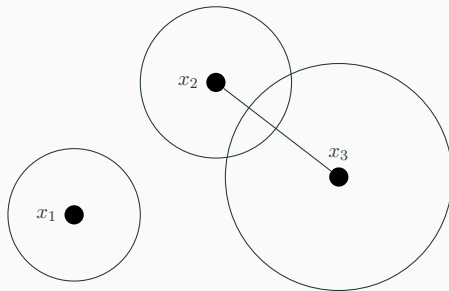
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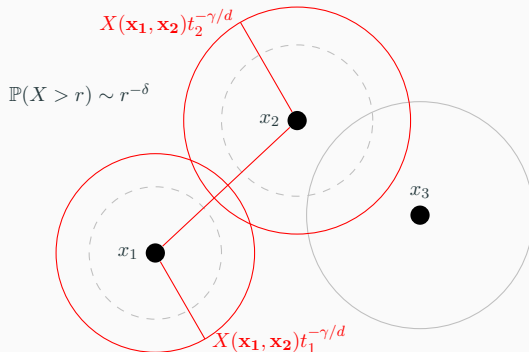
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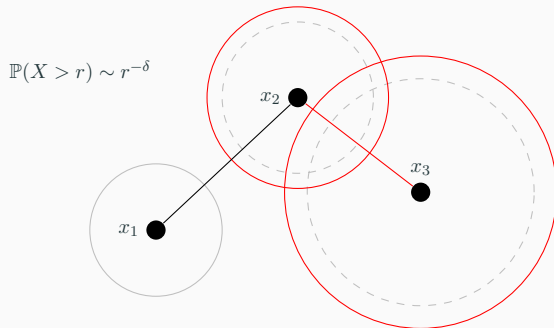
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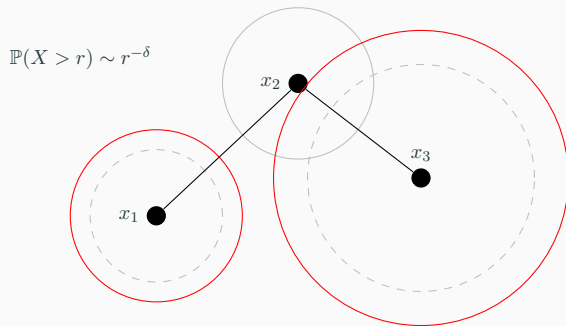
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# **The contact process on the weight-dependent random connection model**

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# The contact process

Let  $G = (V, E)$  be a locally-finite graph.

The **contact process** on  $G$  is a continuous time Markov process  $(\xi_t)_{t \geq 0}$  on  $\{0, 1\}^V$ .

For  $\{x \in V : \xi_t(x) = 1\} \subset V$  we have the transition rates:

$$\begin{aligned} \xi_t &\rightarrow \xi_t \setminus \{x\} && \text{for } x \in \xi_t \text{ at rate } 1, \text{ and} \\ \xi_t &\rightarrow \xi_t \cup \{x\} && \text{for } x \notin \xi_t \text{ at rate } \lambda \cdot |\{y \in \xi_t : x \sim y\}|. \end{aligned}$$

## Key questions

The process has the single absorbing state equal to  $\emptyset$ .

Consequently, the natural questions concern the **extinction time**

$$\tau_G := \inf\{t > 0 : \xi_t = \emptyset\}.$$

- Can  $\tau_G$  be infinite?
- If yes, with what probability?
- On a graph with  $n$  vertices, how does  $\tau_G$  change with  $n$ ?
- How does  $\lambda$  affect the answer?



## Non-extinction probability

Let the contact process with parameter  $\lambda$  on  $G_{(0,T_0)}$  start in the origin  $(0, T_0)$ .

Define:

$$\Gamma(\lambda) = \mathbb{P}_{(0,T_0)}(\xi_t^{(0,T_0)} \neq \emptyset \forall t \geq 0).$$

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### Theorem (G. and Grauer, '24)

*Let  $G$  be a general geometric random graph which satisfies Assumption (A1) for  $\gamma > \frac{\delta}{\delta+1}$ . Then, as  $\lambda \rightarrow 0$ ,*

$$\Gamma(\lambda) \asymp \frac{\lambda^{2/\gamma-1}}{\log(1/\lambda)^{(1-\gamma)/\gamma}}.$$

Some observations:

- The contact process can survive a constant amount of time on a “star” of size at least  $\sim \lambda^{-2}$ . (Mountford, Valesin, and Yao, 2013)
- On a chain of disjoint stars of size at least  $\sim \lambda^{-2} \log(1/\lambda)$ , separated by at most bounded many steps, the contact process survives with probability at least  $p$  for some constant  $p > 0$ . (Linker, Mitsche, Schapira, and Valesin, 2021)

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**Lower bound:** A chain of star  $\rightarrow$  connector  $\rightarrow$  bigger star  $\rightarrow \dots$  exists a.s.

**Upper bound:** Truncated first moment bound counting argument, based on the counting argument by G., *Grauer and Mörters, 2022*.

## Exponential extinction time on finite restrictions

Let  $G_n$  be the restriction of  $G$  to  $\left[-\frac{n^{1/d}}{2}, \frac{n^{1/d}}{2}\right]^d$ .

We denote by  $\tau_n := \inf\{t > 0 : \xi_t^{G_n} = \emptyset\}$  the extinction time of the contact process on  $G_n$ .

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## Theorem (G. and Grauer, '24)

*Let  $(G_n)_{n \in \mathbb{N}}$  be the restricted finite graph sequence of a general geometric random graph which satisfies Assumption (A1) for  $\gamma > \frac{\delta}{\delta+1}$ , and let the graph start fully infected. For any  $\lambda > 0$ , there exists  $c > 0$  such that*

$$\lim_{n \rightarrow \infty} \mathbb{P}\{\tau_n \geq e^{cn}\} = 1.$$

### Proposition (G. and Grauer, '24)

*Let  $S > 0$  be given and  $(G_n)_{n \in \mathbb{N}}$  the restricted finite graph sequence of a general geometric random graph which satisfies Assumption (A1) for  $\gamma > \frac{\delta}{\delta+1}$ . Then, there exists  $b > 0$  and  $\varepsilon > 0$  such that, for  $n$  sufficiently large, the probability that  $G_n$  has a connected subgraph containing  $b \cdot n$  disjoint stars of at least  $S$  vertices each is larger than  $1 - \exp(-n^\varepsilon)$ .*

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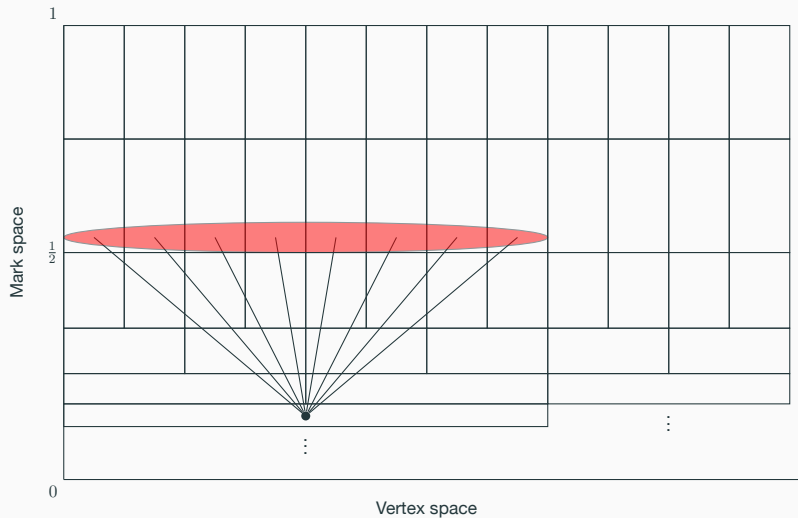
The proof by *Mountford, Mourrat, Valesin, and Yao, 2016* for the configuration model then yields the exponential extinction time claim.



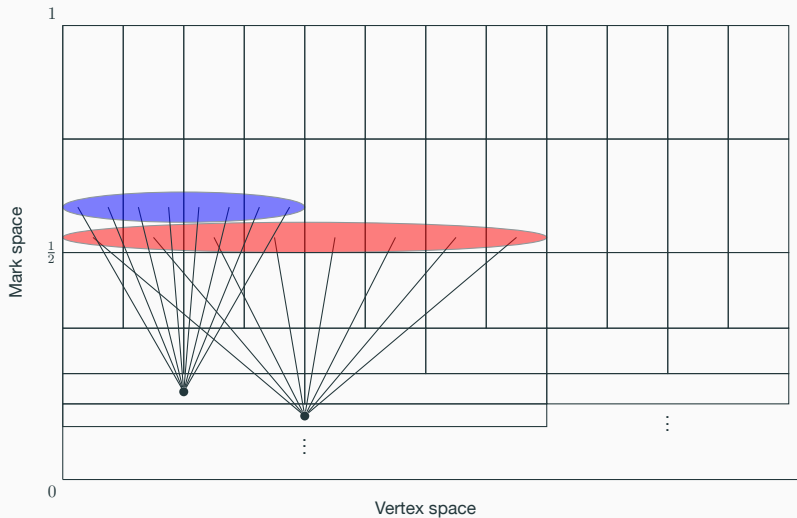
## Proof sketch (Subgraph construction)

The diagram illustrates a 2D grid representing a Markov chain. The vertical axis is labeled "Mark space" with values 0 and  $\frac{1}{2}$ . The horizontal axis is labeled "Vertex space" with values 0 and 1. The grid contains cells labeled  $A_{i,j}$ ,  $B_{i,j}$ ,  $C_{i,j}$ , and  $C_{i,k}$ . The cells are arranged in a 4x12 grid. The first two rows are labeled  $A_{i,j}$  and  $B_{i,j}$ . The third row is labeled  $C_{i,j}$ . The fourth row is labeled  $C_{i,k}$ . The cells are arranged in a 4x12 grid. The first two rows are labeled  $A_{i,j}$  and  $B_{i,j}$ . The third row is labeled  $C_{i,j}$ . The fourth row is labeled  $C_{i,k}$ .

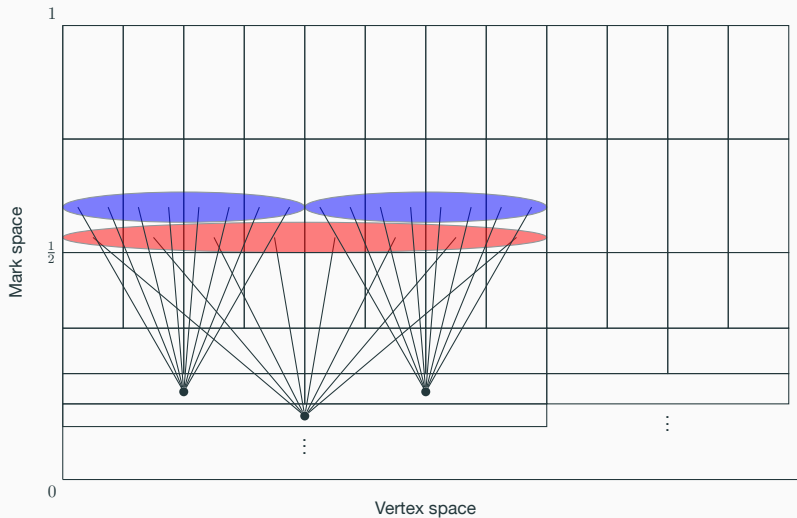
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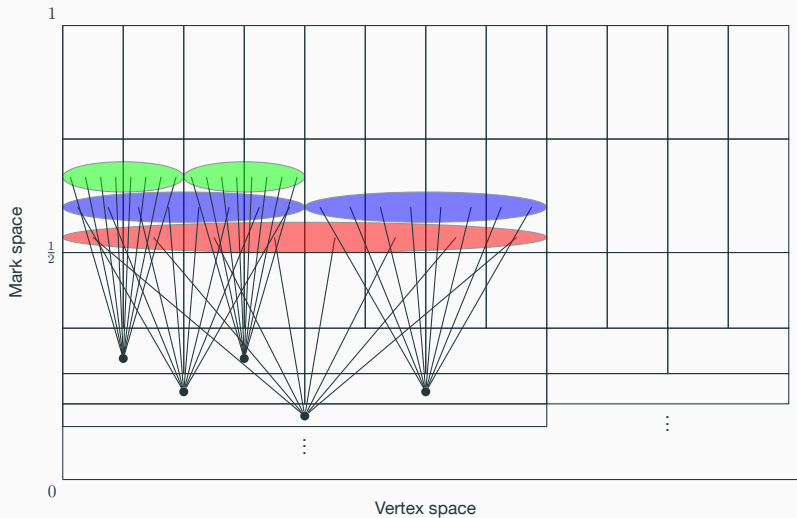
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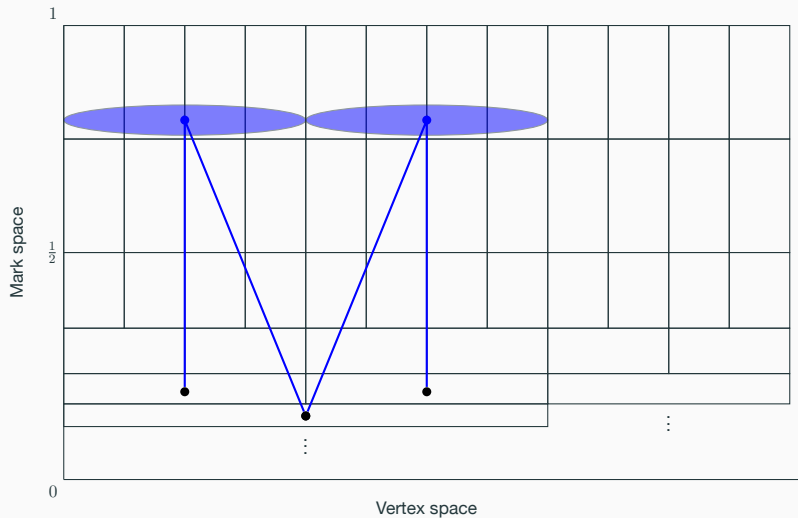
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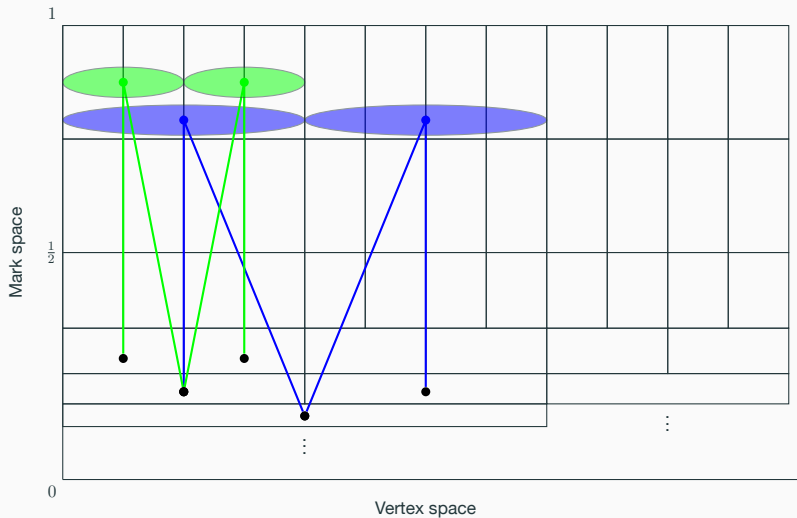
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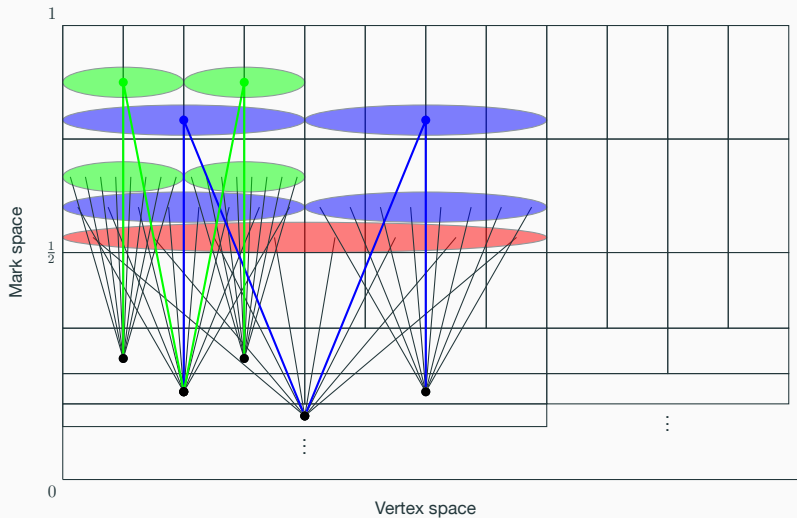
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## Beyond the weight-dependent random connection graph satisfying (A1)

The stated results also hold for the weight-dependent random connection model with the factor kernel, i.e.

$$\mathbb{P}_{\mathbf{x}, \mathbf{y}}(\mathbf{x} \sim \mathbf{y}) \asymp t^{-\delta\gamma} s^{-\delta\gamma} |x - y|^{-\delta d}$$

for  $\delta > 1$  and  $\gamma > \frac{1}{2}$  (Linker, Mitsche, Schapira, and Valesin, 2021).

Key difference: stars connect directly.

## **The dynamic weight-dependent random connection model**

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# Setup

- **Vertices:** a Poisson point process  $\mathcal{X}$  of intensity  $\lambda$  on  $\mathbb{R}^d \times [0, 1]$  or  $\mathbb{T}_n^d \times [0, 1]$  at time  $t = 0$ .

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- **Marks:** Marks do not update\*.

# Setup

- **Edges:** Given  $\mathcal{X}_0$ , an edge is drawn between every pair of vertices  $\mathbf{x} = (x, t), \mathbf{y} = (y, s) \in \mathcal{X}$  if  $U_{x,y} < \mathbb{P}_{\mathbf{x},\mathbf{y}}(\mathbf{x} \sim \mathbf{y})$ , where  $\mathbb{P}_{\mathbf{x},\mathbf{y}}(\mathbf{x} \sim \mathbf{y})$  satisfies (A1), where  $U_{x,y} \sim \text{Unif}(0, 1)$ , sampled independently for every pair  $\{\mathbf{x}, \mathbf{y}\}$ .

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  - This ensures that **conditional** on the locations and marks of two vertices, the existence of an edge between them at times  $t_1$  and  $t_2$  is independent, if  $\lfloor t_1 \rfloor \neq \lfloor t_2 \rfloor$ .
  - The above two events are not unconditionally independent!

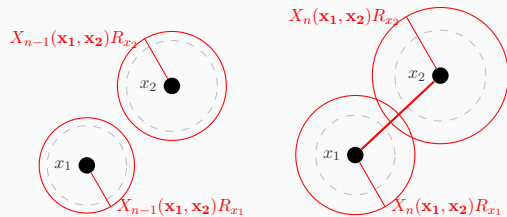


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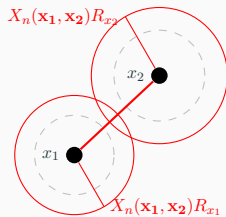
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Individual edge/vertex neighbourhood updating at i.i.d. exponentially distributed random times makes the proofs more difficult, but does not change the results.

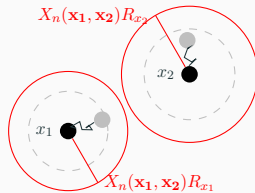
## Edge updating (Example)



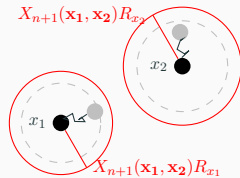
(a) Just before time  $n$ .



(b) At time  $n$ .



(c) Just before  $n + 1$ .



(d) At time  $n + 1$ .

The randomness in the model comes from:

- The **starting locations** and **marks** of the vertices.
- The **motion** of the vertices.
- Given the above, the **random occurrence of the edges**.

## Percolation time

Consider the dynamic scale-free geometric random graph on  $\mathbb{R}^d \times (0, 1)$ .

We define the *percolation time* as

$$T_{\text{perc}} := \inf\{t \geq 0 : \exists \mathbf{x} \in \mathcal{C}_{\infty}^t \text{ s.t. } \mathbf{x} \sim \mathbf{0}\},$$

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### Theorem (G. and Grauer, '24<sup>1</sup>)

Let  $\gamma > \frac{\delta}{\delta+1}$ . Then there exists a constant  $c > 0$  such that on the  $\mathbb{R}^d$  with  $d \geq 1$ , the percolation time  $T_{\text{perc}}$  satisfies

$$\mathbb{P}(T_{\text{perc}} > t) \leq \exp\{-ct^{1/c}\},$$

for any vertex intensity  $\lambda > 0$  and any  $t > 0$  sufficiently large.

<sup>1</sup>Also holds for factor kernel in “robust” phase, i.e.  $\gamma > \frac{1}{2}$ .

## Broadcast time

Let  $\mathbb{T}_n$  be the  $d$ -dimensional torus of volume  $n$  and consider the dynamic scale-free geometric random graph on  $\mathbb{T}_n \times (0, 1)$ .

At time  $t = 0$  an arbitrary vertex starts broadcasting information to all vertices in its connected component. Every vertex that receives this broadcast begins broadcasting it further.

The *broadcast time*  $T_{bc}$  is defined as the smallest time at which every vertex of the network has received the broadcast.

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### **Theorem (G. and Grauer, '24<sup>2</sup>)**

*On the  $d$ -dimensional torus of volume  $n$  with  $d \geq 1$  with  $\gamma > \frac{\delta}{\delta+1}$ , the broadcast time  $T_{bc}$  is with high probability  $O(\log n (\log \log n)^\epsilon)$  for any  $\epsilon > 0$  and any  $\lambda > 0$ .*

<sup>2</sup>Also holds for factor kernel in “robust” phase, i.e.  $\gamma > \frac{1}{2}$ .

## A few remarks

- Dynamic geometric random graphs are strongly correlated spatially and temporally.
- Since the vertex marks are not being updated, exceptionally “unlucky” vertices are possible.
- It is not *a priori* clear how the largest connected component evolves over time.



## Static tools: Good cubes

Let  $a \in (0, \frac{1}{\log 2})$  and  $\Theta \in (\frac{\log 2}{\gamma + \gamma/\delta}, \log 2)$  be two constants.

### Definition ( $t$ - $\alpha$ -dense cubes)

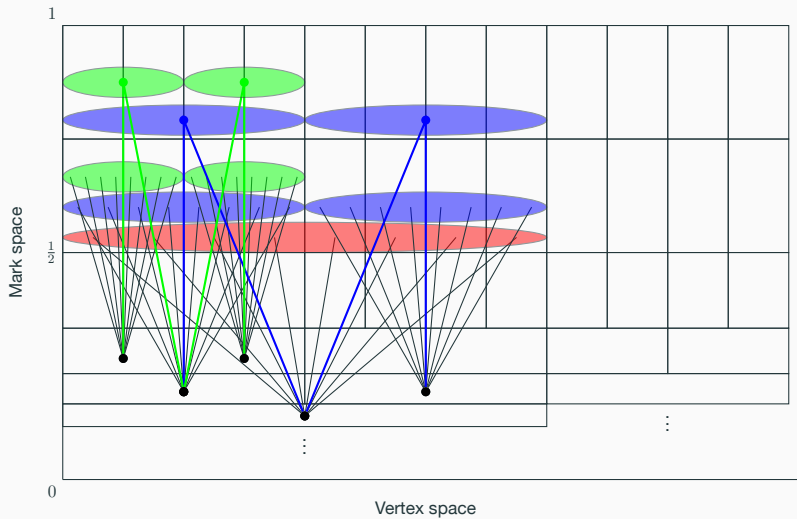
Let  $t > 0$  and  $I_k := (\frac{1}{2}e^{-(k+1)\Theta d}, \frac{1}{2}e^{-k\Theta d})$ ,  $k \in \{0, \dots, \lfloor (a \log t)/d \rfloor\}$  and  $I_{-1} := (\frac{1}{2}, 1)$ . We say a cube  $Q \subset \mathbb{R}^d$  is  $t$ - $\alpha$ -dense, if for every  $k \in \{-1, 0, \dots, \lfloor (a \log t)/d \rfloor\}$ , the locations of the vertices in  $Q \times I_k$  contain as a subset an independent Poisson point process of intensity  $\lambda(1 - \alpha)|I_k|$  on  $\mathbb{R}^d$ , with marks in  $I_k$ .

## Static tools: Evenly spread subgraphs

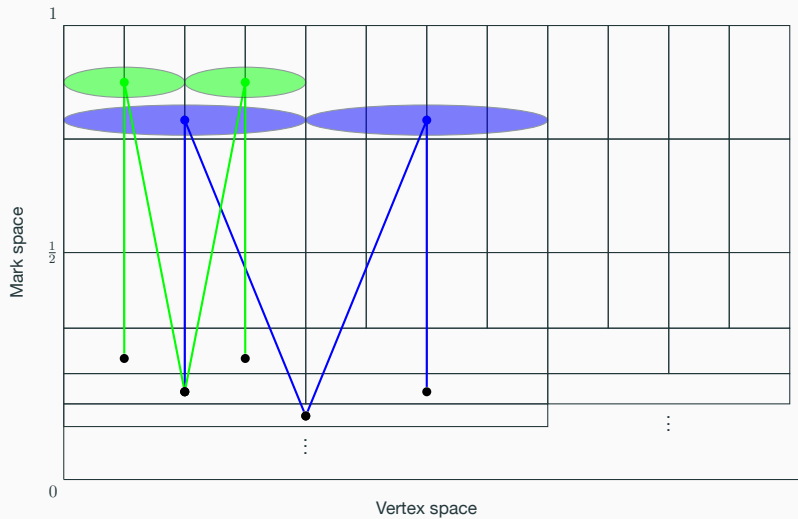
### Definition (Evenly spread subgraphs)

We call a finite connected subgraph of  $\mathcal{G}$  contained inside  $Q_K$  an *evenly spread subgraph of  $\mathcal{G}$  inside  $Q_K$* , if it contains at least  $b \cdot K$  vertices for some constant  $b > 0$  and if every subcube of  $Q_K$  of the form  $\times_{i=1}^d (2^{k_p} v_i, (2^{k_p} + 1) v_i)$ ,  $v_i \in \mathbb{Z}$ ,  $k_p = \lfloor a \log K \rfloor$  contains a vertex with mark smaller than  $\frac{1}{2} e^{-k_p \Theta d}$  belonging to the evenly spread component. We call these vertices the *bottom vertices* of the evenly spread component.

# Evenly spread



# Evenly spread



## Static tools: Evenly spread subgraphs

### Proposition (G. and Grauer, '24)

*Fix  $K$  large enough and consider the cube  $Q_K \subset \mathbb{R}^d$ . Assume that  $Q_K$  is  $K$ - $\alpha$ -dense for some  $\alpha > 0$  and let  $\gamma > \frac{\delta}{\delta+1}$ . Then, there exists  $\epsilon > 0$  such that for any  $\lambda > 0$  and  $K$  large enough, there exists an evenly spread subgraph of  $\mathcal{G}$  inside  $Q_K$  with probability at least*

$$1 - \exp\{-K^\epsilon\}.$$

## Static tools: Evenly spread subgraphs

### Lemma (G. and Grauer, '24)

*Assume that  $Q_K$  is  $K$ - $\alpha$ -dense for some  $\alpha > 0$  and let  $\gamma > \frac{\delta}{\delta+1}$  and  $\lambda > 0$ . Let furthermore  $G$  be an evenly spread subgraph of  $\mathcal{G}^{1-\varepsilon}$  inside  $Q_K$ . Then, a given vertex  $x$  with a mark in  $(0, 1)$  and arbitrary location within  $Q_K$  belongs to the same connected component of  $\mathcal{G}$  as  $G$  with probability bounded away from 0.*

## Static tools: Summary

- If a cube  $Q_K$  of volume  $t^d$  is  $t$ - $\alpha$ -dense, an evenly spread subgraph of size  $K$  exists  $wep(K)$ .
- An arbitrary vertex anywhere in the cube  $Q_K$  has probability bounded away from 0 of belonging to the same connected component as the evenly spread subgraph.

### Key question

How often is the cube  $Q_K$   $t$ - $\alpha$ -dense during the time interval  $[0, t]$ ?

## Proposition (Peres, Sinclair, Sousi, and Stauffer, 2011<sup>3</sup>)

Fix  $K > \ell > 0$  and consider the cube  $Q_{K^d} \subset \mathbb{R}^d$  tessellated into subcubes of side length  $\ell$ . Let  $\Pi_0$  be an arbitrary point process at time 0 that contains at least  $\beta \ell^d$  vertices in each subcube of the tessellation for some  $\beta > 0$ . Let  $\Pi_\Delta$  be the point process obtained at time  $\Delta$  from  $\Pi_0$  after the vertices have moved for time  $\Delta$ . Fix  $\epsilon \in (0, 1)$  and let  $\Psi$  be an independent Poisson point process of intensity  $(1 - \epsilon)\beta$  on  $Q_K$ . Then there exists a coupling of  $\Psi$  and  $\Pi_\Delta$  and constants  $c_1, c_2, c_3$  that depend on  $d$  only, such that if  $\Delta \geq \frac{c_1 \ell^2}{\epsilon^2}$  and  $K' \leq K - c_2 \sqrt{\Delta \log \epsilon^{-1}} > 0$ , the vertices of  $\Psi$  are a subset of the vertices  $\Pi_\Delta$  inside the cube  $Q_{K'}$  with probability at least

$$1 - \frac{K^d}{\ell^d} \exp\{-c_3 \epsilon^2 \beta \Delta^{d/2}\}.$$

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<sup>3</sup>Generalisations of this result exist also for (uniformly elliptic) lattices and fractal lattices.



## Dynamics: Cubes are good/dense most of the time

### Proposition (G. and Grauer, '24)

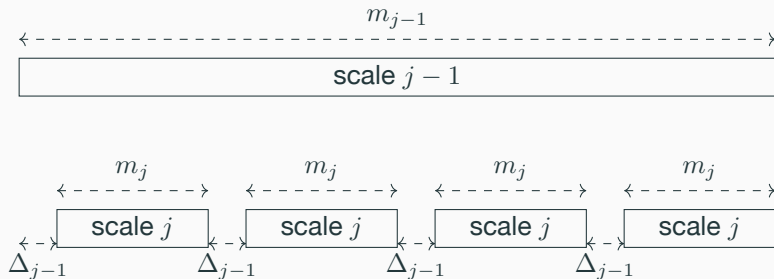
Let  $t > 0$  be a sufficiently large integer and  $\xi, \epsilon \in (0, 1)$  two constants. Consider the cube  $Q_{L^d}$ , for  $L = t$ . Define for  $i = 0, \dots, t$  the events

$$A_i = \{\text{at time } i \text{ the cube } Q_{L^d} \text{ is } t\text{-}\xi\text{-dense}\}.$$

Then, there exist two positive constants  $c_1, c_2$  such that

$$\mathbb{P} \left( \sum_{i=0}^{t-1} \mathbb{1}_{A_i} \geq (1 - \epsilon)t \right) \geq 1 - \exp\{-c_1 t^{c_2}\}.$$

## Dynamics: Proof of proposition (sketch)

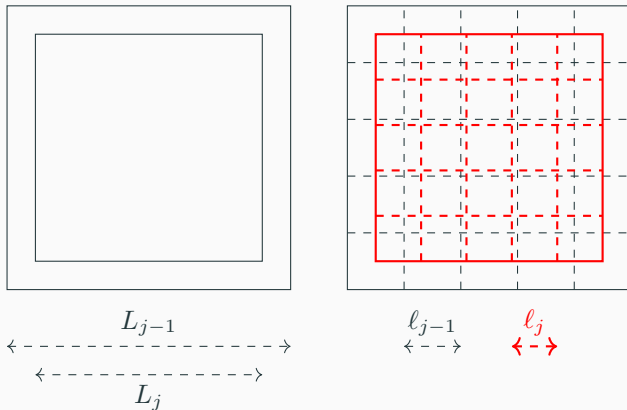


*The temporal multi-scale recursion*

The values are chosen as:

$$m_1 = t, \quad m_j = \frac{m_{j-1} - 4\Delta_{j-1}}{4}, \quad \text{and} \quad \frac{\Delta_j}{m_j} = \frac{\epsilon}{8\kappa}.$$

## Dynamics: Proof of proposition (sketch)



*The spatial multi-scale recursion*

The values are chosen as:  $L_1 = L^2$ ,  $L_\kappa = L$ ,  $\ell_1 = c\sqrt{t}$ ,  $\ell_{j+1} = \ell_j \sqrt{\frac{1}{4} - \frac{\epsilon}{8\kappa}}$ , with  $\kappa = O(\log t)$ .

## Dynamics: Proof of proposition (sketch)

Let  $D_j$  be the event that all subcubes of side-length  $\ell_j$  inside the cube of side length  $L_j$  are *good* for a fraction of at least  $(1 - \frac{\epsilon_j}{2})$  time intervals of scale  $j$ .

1.  $\mathbb{P}(D_1) \geq 1 - \exp\{-c_1 t^{c_2}\}$

2. Set

$$E = \{\text{at time } b' \text{ all subcubes are good for the scale } j-1\},$$

with  $b' = b - \Delta_{j-1}$ . Then

$$\mathbb{P}([b, b + m_j] \text{ is good}, E|F) \geq 1 - \exp\{-c_1 t^{-a\Theta} \ell_j^d / \kappa^2\},$$

with  $F$  measurable w.r.t. events that occurred up to time  $b'$ .

3.  $\mathbb{P}(D_j^c \cap D_{j-1}) \leq \exp\{-\frac{ct^{c_4}}{\kappa^6}\}.$

## Dynamics: Proof of proposition (sketch)

Consequently,

$$\mathbb{P}(D_{\kappa}^c) \leq \mathbb{P}(D_{\kappa}^c \cap D_{\kappa-1}) + \mathbb{P}(D_{\kappa-1}^c),$$

which gives

$$\mathbb{P}(D_{\kappa}^c) \leq \sum_{j=2}^{\kappa} \mathbb{P}(D_j^c \cap D_{j-1}) + \mathbb{P}(D_1^c).$$

## Broadcast time

Consider a torus of volume  $n$  and let  $t = C \log n (\log \log n)^\epsilon$ .

- Tessellate the torus into cubes of side length  $t$ .
- Each such cube is dense at least  $(1 - \epsilon)t$  amount of time  $\Rightarrow$  Contains  $wep(t)$  evenly spread subgraph throughout this time.
- At each time, the second largest component on the torus (of size  $n$ ) is  $wep(t)$  of size  $o(t)$  by *Jorritsma, Komjáthy, and Mitsche, 2024*.

Consequently, when an evenly spread subgraph exists, it is in the largest connected component of the torus.

## Broadcast time

- Each vertex (including the “origin vertex”) has positive (bdd. away from 0) probability of belonging to the same component as the “local” evenly spread subgraph.  
⇒ Infection enters the large component during  $[0, t/2)$  with probability greater than  $1 - (1 - p)^{t/2}$ .
- During  $[t/2, t]$ , a vertex belongs to the large component at least once with probability greater than  $1 - (1 - p)^{t/2}$ .
- Since there are only  $\Theta(n)$  vertices on the torus  $whp(n)$ , every vertex gets the information with probability at least

$$1 - (1 + \delta)n(1 - p)^{(1-\epsilon)t/2}.$$


Setting  $C$  large concludes the proof.




## Open problems and future work

- Tightness of broadcast time bound.
- **Non-robust regimes.**
- Different edge updating mechanisms.
- **Contact process on dynamic graph.**



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**Thank you for listening!<sup>4</sup>**

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<sup>4</sup>and/or being physically present.