

The differential equation method

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Describing elements of a random graph

Consider a graph G on n vertices where edges are present independently at random

Vertex set $[n] = \{1, 2, \dots, n\}$

Edges r.v.'s $X_{ij} = \mathbb{1}_{i \leftrightarrow j}$ independent

Some interesting statistics:

- number of edges
- isolated vertices
- number of components (of given size)
- degree sequence
- chromatic number
- independence number

Intuitively changing one variable in $\{X_{ij}\}$ won't change statistics by much.

This course: conditions to 'see' trends ... in a (graph) discrete process

Chernoff bounds $t > 0$

$$\mathbb{P}(X \geq a) = \mathbb{P}(e^{tX} \geq e^{ta}) \leq \frac{\mathbb{E}[e^{tX}]}{e^{ta}}$$

$$t = 4a / \sum (b_i - a_i)^2 \quad \text{optimize } t$$

$$\text{If } S_n = \sum_{i=1}^n X_i \text{ independent } X_i$$

$$\mathbb{P}(S_n \geq \mathbb{E}[S_n] + \varepsilon) \leq e^{-t\varepsilon} \prod_{i=1}^n \mathbb{E}[e^{t(X_i - \mathbb{E}[X_i])}]$$

+ bounded X_i so that $a_i \leq X_i \leq b_i$

$$\mathbb{P}(S_n \geq \mathbb{E}[S_n] + \varepsilon) \leq e^{-2\varepsilon^2 / \sum (b_i - a_i)^2}$$

This gives quantitative bounds for LLN $\mathbb{P}(X_i \in (0, 1))$

$$\sum_{n \geq 1} \mathbb{P}\left(\left|\frac{S_n - \mathbb{E}[S_n]}{n}\right| > u\right) \leq \sum_{n \geq 1} 2e^{-nu^2/2} < \infty$$

$$\Rightarrow_{\text{B.C.}} \mathbb{P}\left(\left|\frac{S_n - \mathbb{E}[S_n]}{n}\right| > u \text{ infinitely often}\right) = 0$$

$$\frac{S_n - \mathbb{E}[S_n]}{n} \xrightarrow{\text{a.s.}} 0$$

Two generalization

① If $\{Z_i\}_{i \geq 0}$ martingale, $Z_0 = 0$

$$|Z_{i+1} - Z_i| \leq c_i$$

$$\mathbb{P}(Z_n \geq \mathbb{E}[Z_n] + \varepsilon) \leq e^{-2\varepsilon^2 / \sum c_i^2}$$

Proof Sketch: $\mathcal{F}_i = \sigma\text{-algebra}(Z_1, \dots, Z_i)$
 $Z_n = \sum_{i=1}^n Z_i - Z_{i-1} = \sum_{i=1}^n V_i$

$$\begin{aligned}\mathbb{E}[e^{tZ_n}] &= \mathbb{E}\left[\prod_{i=1}^n e^{tV_i}\right] \\ &= \mathbb{E}\left[\prod_{i=1}^n e^{tV_i} \mathbb{E}[e^{tV_i} | \mathcal{F}_{i-1}]\right]\end{aligned}$$

Now V_n satisfies conditions of Hoeffding's lemma.

② Let $\{X_i\}$ independent $X_i \in \mathcal{X}$

$$f: \mathcal{X}^n \rightarrow \mathbb{R}$$

Conditions

$$|f(x_1, x_2, \dots, x_n) - f(y_1, \dots, y_n)| \leq \sum c_i \mathbb{1}_{x_i \neq y_i}$$

$$\mathbb{P}(f(X) \geq \mathbb{E}[f] + \varepsilon) \leq e^{-2\varepsilon^2 / \sum c_i^2}$$

$$V_i = \mathbb{E}[f | \mathcal{F}_i] - \mathbb{E}[f | \mathcal{F}_{i-1}]$$

$$\mathcal{F}_i = \sigma\text{-algebra}(X_1, \dots, X_i)$$

$\exists$ r.v.'s L_i, U_i measurable w.r.t $\mathcal{F}_{i-1}$

$$L_i \leq V_i \leq U_i \quad U_i - L_i \leq c_i$$

so can apply Hoeffding's lemma

Extension in Hamming distance

$$\underline{x} = (x_1, \dots, x_n) \in S^n \quad A \subset S^n$$

Hamming dist $d_H(\underline{x}, y) = \sum \mathbb{1}_{x_i \neq y_i}$

$$d_H(\underline{x}, A) = \inf_{y \in A} d_H(\underline{x}, y)$$

$$[A]_r = \{ \underline{x} \in S^n : d_H(\underline{x}, A) \leq r \}$$

How big becomes $[A]_{\epsilon n}$

$$\begin{aligned} \mathbb{P}(\underline{x} \in [A]_{\epsilon n}) &= \mathbb{P}(d(\underline{x}, A) \leq \epsilon n) \\ &= 1 - \mathbb{P}(d(\underline{x}, A) > \epsilon n) \end{aligned}$$

$$* \epsilon n > \mathbb{E}[d(\underline{x}, A)] \quad u = \epsilon n - \mathbb{E}[d(\underline{x}, A)]$$

$$\mathbb{P}(d(\underline{x}, A) > \epsilon n) = \mathbb{P}(d(\underline{x}, A) > \mathbb{E}[\cdot] + u) \leq e^{-2u^2}$$

Problem: $\mathbb{E}[d(\underline{x}, A)]$ unknown

$$\epsilon n \geq B \geq \mathbb{E}[d(\underline{x}, A)]$$

$$\mathbb{P}(d(\underline{x}, A) \geq \epsilon n) \leq \mathbb{P}(d(\underline{x}, A) \geq \mathbb{E}[\cdot] + u)$$

$$u = \epsilon n - B$$

Trick: from Hoeffding lemma 'proof

$$\mathbb{E}[e^{t(f - \mathbb{E}[f])}] \leq e^{nt^2/8}$$

$$f = -d(\underline{x}, A) \quad \mathbb{E}[e^{t\mathbb{E}[f]}] = e^{-t\mathbb{E}[f]}$$

$$\mathbb{P}(A) = \mathbb{E}[\mathbb{1}_A e^{tf}] \leq \mathbb{E}[e^{tf}] = \mathbb{E}[e^{-td(\underline{x}, A)}]$$

$$\mathbb{P}(A) e^{-t\mathbb{E}[d(\underline{x}, A)]} \leq e^{nt^2/8}$$

Take logarithm + optimize t

$$\mathbb{E}[d(\underline{x}, A)] \leq \sqrt{\frac{-n \log \mathbb{P}(A)}{2}} = B$$