

Previously in MA30087/50087:

- transporting goods from **sources** S_1, \dots, S_m with supplies s_1, \dots, s_m to **destinations** D_1, \dots, D_n with demands d_1, \dots, d_n
- **amount transported** from S_i to D_j is x_{ij}
- **cost of transporting** a unit from S_i to D_j is c_{ij}
- **balanced transportation problem** (so $\sum_{i=1}^m s_i = \sum_{j=1}^n d_j$)

$$\begin{aligned} \text{minimise} \quad & z = \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \\ \text{subject to} \quad & \sum_{j=1}^n x_{ij} = s_i \quad i = 1, \dots, m \\ & \sum_{i=1}^m x_{ij} = d_j \quad j = 1, \dots, n \\ & x_{ij} \geq 0 \quad i = 1, \dots, m, j = 1, \dots, n \end{aligned}$$

Today in MA30087/50087:

$$\mathbf{Ax} = \mathbf{b} \Rightarrow \begin{pmatrix} 1 & \cdots & 1 & \cdots & 0 & \cdots & 0 \\ \vdots & \ddots & \vdots & \ddots & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & \cdots & 1 & \cdots & 1 \\ 1 & \cdots & 0 & \cdots & 1 & \cdots & 0 \\ \vdots & \ddots & \vdots & \ddots & \vdots & \ddots & \vdots \\ 0 & \cdots & 1 & \cdots & 0 & \cdots & 1 \end{pmatrix} \begin{pmatrix} x_{11} \\ \vdots \\ x_{1n} \\ \vdots \\ x_{m1} \\ \vdots \\ x_{mn} \end{pmatrix} = \begin{pmatrix} s_1 \\ \vdots \\ s_m \\ d_1 \\ \vdots \\ d_n \end{pmatrix}$$

- **dual of the balanced transportation problem**, $\mathbf{y} = (u_1, \dots, u_m, v_1, \dots, v_n)^T$

$$\begin{aligned} \text{maximise} \quad & z' = \sum_{i=1}^m u_i s_i + \sum_{j=1}^n v_j d_j \\ \text{subject to} \quad & u_i + v_j \leq c_{ij} \quad i = 1, \dots, m, j = 1, \dots, n \end{aligned}$$

- balanced transportation problem has a finite optimal solution
- $\text{Rank}(\mathbf{A}) = m + n - 1$
- transportation algorithm utilises **asymmetric complementary slackness**

$$x_{ij}[u_i + v_j - c_{ij}] = 0 \quad \forall i = 1, \dots, m, j = 1, \dots, n$$