

19/10/2015.

LECTURE 8: CBI.11.

(A)

LAST TIME: $\underline{x} = \begin{pmatrix} \underline{x}_B \\ \underline{0}_{n-m} \end{pmatrix}, \underline{v} = \begin{pmatrix} \underline{v}_B \\ \underline{0}_{n-m} \end{pmatrix}, \underline{w} = \begin{pmatrix} \underline{w}_B \\ \underline{0}_{n-m} \end{pmatrix}$

NOW, $\underline{b} = \underline{A} \underline{v} = (\underline{B} | \underline{N}) \begin{pmatrix} \underline{v}_B \\ \underline{0}_{n-m} \end{pmatrix} = \underline{B} \underline{v}_B$

THUS, AS $\underline{x}, \underline{v} \in F_{(1)}$, $\underline{b} = \underline{B} \underline{x}_B = \underline{B} \underline{v}_B$ SO THAT

$$\underline{B} (\underline{x}_B - \underline{v}_B) = \underline{0}_m \text{ i.e. } \sum_{j=1}^m (x_j - v_j) \underline{a}_j = \underline{0}_m.$$

BUT $\underline{a}_1, \dots, \underline{a}_m$ ARE LINEARLY INDEPENDENT SO $x_j - v_j = 0$
 $\forall j = 1, \dots, m.$

i.e. $\underline{x}_B = \underline{v}_B \Rightarrow \underline{x} = \underline{v}$ AND HENCE $\underline{v} = \underline{w}$

THUS, $\underline{v}, \underline{w}$ ARE NOT DISTINCT: CONTRADICTION. □

THEOREM.

CONSIDER A LP PROBLEM IN CANONICAL FORM. SUPPOSE $\underline{x}$ IS AN EXTREME POINT OF $F_{(1)}$ THEN $\underline{x}$ IS A BASIC FEASIBLE SOLUTION.

PROOF.

WE NEED TO SHOW THAT THE COLUMNS OF $\underline{A}$ FOR WHICH $x_j > 0$ ARE LINEARLY INDEPENDENT.

WLOG ASSUME FIRST k COMPONENTS OF $\underline{x}$ ARE NON-ZERO.

$$\text{i.e. } x_j > 0 \text{ } j=1, \dots, k \quad x_j = 0 \text{ } j=k+1, \dots, n.$$

$$\underline{A} \underline{x} = \sum_{j=1}^n x_j \underline{a}_j = \sum_{j=1}^k x_j \underline{a}_j = \underline{b} \quad (*)$$

** ASSUME THAT $\underline{a}_1, \dots, \underline{a}_k$ ARE LINEARLY DEPENDENT **

(B).

[AIM: PROOF BY CONTRADICTION. WE'RE GOING TO FIND TWO POINTS IN F_C , WHICH $\underline{x}$ LIES ON A SEGMENT BETWEEN SO THAT $\underline{x}$ IS NOT EXTREME].

IN WHICH CASE, THERE EXIST CONSTANTS α_j NOT ALL EQUAL TO ZERO SUCH THAT

$$\sum_{j=1}^k \alpha_j \underline{a}_j = \underline{0}_m \quad (**)$$

[ADDING MULTIPLES OF (**) TO (*) CAN SATISFY $A\underline{x}' = \underline{b}$. WANT TO MAINTAIN $\underline{x}' \geq \underline{0}_n$].

LET $\mu > 0$ BE SUCH THAT $\mu |\alpha_j| \leq x_j$ FOR ALL $j=1, \dots, k$.
HENCE,

$$x_j - \mu \alpha_j \geq x_j - \mu |\alpha_j| \geq 0 \quad \begin{cases} x_j - \mu |\alpha_j| & \text{if } \alpha_j > 0 \\ x_j + \mu |\alpha_j| & \text{if } \alpha_j < 0 \end{cases}$$

$$x_j + \mu \alpha_j \geq x_j - \mu |\alpha_j| \geq 0 \quad \begin{cases} x_j + \mu |\alpha_j| & \text{if } \alpha_j > 0 \\ x_j - \mu |\alpha_j| & \text{if } \alpha_j < 0 \end{cases}$$

FROM (*) AND (**)

$$\sum_{j=1}^k (x_j - \mu \alpha_j) \underline{a}_j = \underline{b}, \quad \sum_{j=1}^k (x_j + \mu \alpha_j) \underline{a}_j = \underline{b}.$$

NOTE: $\underline{\alpha} \neq \underline{0}_n$

LET $\underline{\alpha} = (\alpha_1, \dots, \alpha_k, 0, \dots, 0)$ THEN $\underline{x} - \mu \underline{\alpha} \in F_C$ AND $\underline{x} + \mu \underline{\alpha} \in F_C$. AS F_C IS CONVEX THEN FOR ALL $\lambda \in [0, 1]$

$$\lambda(\underline{x} - \mu \underline{\alpha}) + (1 - \lambda)(\underline{x} + \mu \underline{\alpha}) \in F_C$$

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TAKING $\lambda = 1/2$ GIVES $\underline{x} = \frac{1}{2} (\underline{x} - \mu \underline{\alpha}) + \frac{1}{2} (\underline{x} + \mu \underline{\alpha})$

SO THAT $\underline{x}$ IS NOT AN EXTREME POINT: A CONTRADICTION.

[AS NOT ALL $\alpha_j = 0$ IN DEPENDENCE ASSUMPTION, $\mu > 0$]
AND $\underline{\alpha} \neq \underline{0}$ □

NOTES.

1. IF $k < m$ THEN $\underline{x}$ IS A DEGENERATE BASIC SOLUTION.
2. WE CAN FIND $m - k$ COLUMNS FROM $\underline{a}_{k+1}, \dots, \underline{a}_n$ SUCH THAT THESE AND $\underline{a}_1, \dots, \underline{a}_k$ FORM A LINEARLY INDEPENDENT SET: SEE Q3 OF QUESTION SHEET THREE

[DEGENERATE EXAMPLE THIS MORNING, $k=1$, $\{\underline{a}_1\}$ $m-2$ SELECT EITHER $\underline{a}_2$ OR $\underline{a}_3$ TO FORM LINEARLY INDEPENDENT SET]

3.2 THE FUNDAMENTAL THEOREM OF LINEAR PROGRAMMING.

THEOREM.

CONSIDER ALP PROBLEM IN CANONICAL FORM WITH FEASIBLE REGION $F(c)$.

- ①. IF THERE IS A FEASIBLE SOLUTION, THERE IS A BASIC FEASIBLE SOLUTION.
- ②. IF THERE IS AN OPTIMAL SOLUTION, THERE IS AN OPTIMAL BASIC FEASIBLE SOLUTION.

PROOF.

WE SHALL PROVE THE RELATIONSHIP BETWEEN OPTIMAL SOLUTIONS; FEASIBLE SOLUTIONS FOLLOW IN A SIMILAR WAY.

①.

LET $\underline{x}$ BE AN OPTIMAL FEASIBLE SOLUTION. WLOG LET NON-ZERO COMPONENTS OF $\underline{x}$ BE THE FIRST k COMPONENTS

$$\text{i.e. } x_j = \begin{cases} > 0 & j=1, \dots, k \\ 0 & j=k+1, \dots, n. \end{cases}$$

WE HAVE $\sum_{j=1}^k x_j \underline{a}_j = \underline{b}$. IF $\underline{a}_1, \dots, \underline{a}_k$ ARE LINEARLY

INDEPENDENT WE'RE DONE. OTHERWISE, SUPPOSE $\underline{a}_1, \dots, \underline{a}_k$ LINEARLY DEPENDENT. THEN THERE EXIST CONSTANTS α_j NOT ALL EQUAL TO ZERO SUCH THAT

$$\sum_{j=1}^k \alpha_j \underline{a}_j = \underline{0}_m$$

[WE'RE GOING TO SHOW THAT AS LONG AS THE $\underline{a}_1, \dots, \underline{a}_k$ ARE LINEARLY DEPENDENT WE CAN FIND AN OPTIMAL SOLUTION THAT HAS ONE OF THE $x_1, \dots, x_k$ EQUAL TO ZERO AND CONTINUE UNTIL WE HAVE INDEPENDENCE]

LET $\mu > 0$ BE SUCH THAT $\mu |\alpha_j| \leq x_j$ FOR ALL $j=1, \dots, k$ THEN (AS IN THE PREVIOUS PROOF)

$$\underline{x} - \mu \underline{\alpha}, \quad \underline{x} + \mu \underline{\alpha} \in F_C \text{ WHERE } \underline{\alpha} = (\alpha_1, \dots, \alpha_k, 0, \dots, 0)^T$$

AS $\underline{x}$ OPTIMAL

$$\underline{c}^T \underline{x} \geq \underline{c}^T (\underline{x} - \mu \underline{\alpha}) = \underline{c}^T \underline{x} - \mu \underline{c}^T \underline{\alpha}$$

$$\text{i.e. } \underline{\mu \underline{c}^T \underline{\alpha}} \geq 0. \quad (*)$$