

Topics: Linda

Linda: an adaption of the thread pool/worker idea

Topics: Linda

Linda: an adaption of the thread pool/worker idea

Like these, it is task based, with threads choosing tasks and executing them

Topics: Linda

Linda: an adaption of the thread pool/worker idea

Like these, it is task based, with threads choosing tasks and executing them

However, now, the tasks have extra structure to guide the choice

Topics: Linda

Linda: an adaption of the thread pool/worker idea

Like these, it is task based, with threads choosing tasks and executing them

However, now, the tasks have extra structure to guide the choice

And the view of the system is flipped from thinking of it as a thread pool to thinking of it as a task pool

Linda

The world is based around *tuples*: these are simple (short) ordered sequences of values

Linda

The world is based around *tuples*: these are simple (short) ordered sequences of values

E.g., [1, "hello"] and [2, "goodbye"]

Linda

The world is based around *tuples*: these are simple (short) ordered sequences of values

E.g., [1, "hello"] and [2, "goodbye"]

And there is a global pool or *tuplespace* containing these tuples

Linda

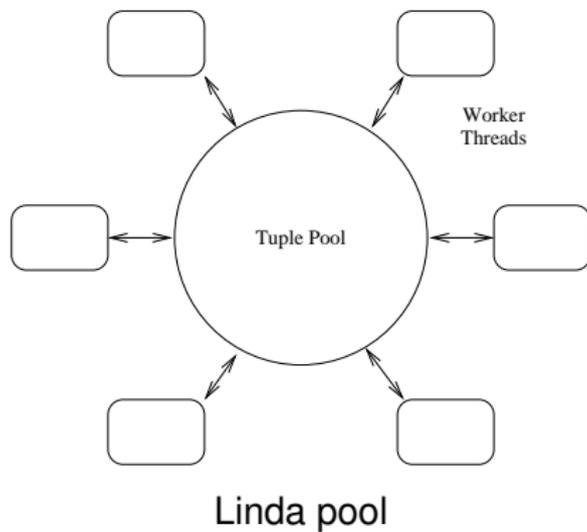
The world is based around *tuples*: these are simple (short) ordered sequences of values

E.g., [1, "hello"] and [2, "goodbye"]

And there is a global pool or *tuplespace* containing these tuples

Threads communicate via the pool by putting tuples in and taking tuples out

Linda



All communications via the pool

Linda

There are four operations the threads can execute:

Linda

There are four operations the threads can execute:

- out send a tuple to the pool

Linda

There are four operations the threads can execute:

- `out` send a tuple to the pool
- `in` get and remove a tuple from the pool

Linda

There are four operations the threads can execute:

- `out` send a tuple to the pool
- `in` get and remove a tuple from the pool
- `read` get but don't remove a tuple from the pool

Linda

There are four operations the threads can execute:

- `out` send a tuple to the pool
- `in` get and remove a tuple from the pool
- `read` get but don't remove a tuple from the pool
- `eval` create a new thread

Linda

The important bit is how in and read work

Linda

The important bit is how `in` and `read` work

The arguments to these are either (a) a literal constant (e.g., string, integer) or (b) a pattern variable, e.g., `?s`, or something suitable for the language you are using

Linda

The important bit is how `in` and `read` work

The arguments to these are either (a) a literal constant (e.g., string, integer) or (b) a pattern variable, e.g., `?s`, or something suitable for the language you are using

A matching tuple is returned from the pool where a match is defined by

- the tuple is the same length and
- the constant literals match

Linda

The important bit is how `in` and `read` work

The arguments to these are either (a) a literal constant (e.g., string, integer) or (b) a pattern variable, e.g., `?s`, or something suitable for the language you are using

A matching tuple is returned from the pool where a match is defined by

- the tuple is the same length and
- the constant literals match

Then the pattern variables are set to the corresponding values in the chosen tuple

Linda

For example, if the pool contains

```
[1, "hello"], [2, "goodbye"], [1, "world"]
```

then there are two matches for `[1, ?s]`, namely `[1, "hello"]` and `[1, "world"]`

Linda

For example, if the pool contains

```
[1, "hello"], [2, "goodbye"], [1, "world"]
```

then there are two matches for `[1, ?s]`, namely `[1, "hello"]` and `[1, "world"]`

One of these will be chosen, *non-deterministically*

Linda

For example, if the pool contains

```
[1, "hello"], [2, "goodbye"], [1, "world"]
```

then there are two matches for `[1, ?s]`, namely `[1, "hello"]` and `[1, "world"]`

One of these will be chosen, *non-deterministically*

If the former is chosen, then the variable `s` will be given the value `"hello"`

Linda

Note that the matching and choosing done by the pool: a possible bottleneck if the implementation of the Linda library is not careful

Linda

Note that the matching and choosing done by the pool: a possible bottleneck if the implementation of the Linda library is not careful

Of course, the pool might itself be a multithreaded system

Linda

Note that the matching and choosing done by the pool: a possible bottleneck if the implementation of the Linda library is not careful

Of course, the pool might itself be a multithreaded system

If no matching tuple exists, the call will block until one arrives

Linda

Note that the matching and choosing done by the pool: a possible bottleneck if the implementation of the Linda library is not careful

Of course, the pool might itself be a multithreaded system

If no matching tuple exists, the call will block until one arrives

If more than one thread simultaneously matches a tuple using `in`, exactly one will get the tuple

Linda

Note that the matching and choosing done by the pool: a possible bottleneck if the implementation of the Linda library is not careful

Of course, the pool might itself be a multithreaded system

If no matching tuple exists, the call will block until one arrives

If more than one thread simultaneously matches a tuple using `in`, exactly one will get the tuple

The action of match and removal for an `in` is atomic

Linda

If both threads use `read`, there is no problem, they both get a copy

Linda

If both threads use `read`, there is no problem, they both get a copy

If one uses an `in` and the other a `read`, it can go either way:

- `read` before `in`: they both get the tuple
- `in` before `read`: the `in` gets the tuple, the `read` doesn't

Linda

If both threads use `read`, there is no problem, they both get a copy

If one uses an `in` and the other a `read`, it can go either way:

- `read` before `in`: they both get the tuple
- `in` before `read`: the `in` gets the tuple, the `read` doesn't

This non-deterministic outcome would normally be considered a programmer error

Linda

These subtleties mean you must be quite careful with Linda

Linda

These subtleties mean you must be quite careful with Linda

A common paradigm is to use an initial *tag*, often an integer, as in `[1, "hello"]`, to impose some structure on the tuples

Linda

Dining Philosophers in Linda

Linda

Dining Philosophers in Linda

We have five philosophers and shall prevent deadlock by only letting four sit at a time

Linda

Dining Philosophers in Linda

We have five philosophers and shall prevent deadlock by only letting four sit at a time

Initial conditions:

```
out("place ticket") four times;  
out("chopstick", i) for  $i = 0 \dots 4$   
eval(phil, i) for  $i = 0 \dots 4$ 
```

Linda

```
defun phil(i) {  
  while true {  
    think()  
    in("place ticket")  
    in("chopstick", i)  
    in("chopstick", i+1 mod 5)  
    eat()  
    out("chopstick", i)  
    out("chopstick", i+1 mod 5)  
    out("place ticket")  
  }  
}
```

This example contains no patterns, only constant literals

Linda

```
defun phil(i) {
  while true {
    think()
    in("place ticket")
    in("chopstick", i)
    in("chopstick", i+1 mod 5)
    eat()
    out("chopstick", i)
    out("chopstick", i+1 mod 5)
    out("place ticket")
  }
}
```

This example contains no patterns, only constant literals

Note Linda does not eliminate the possibility of deadlock in badly written programs: just put the (in "place ticket") after the in of the chopsticks

Linda

Producers/Consumers is just as easy

```
defun producer(n) {
  out(n, make-product())
  producer(n + 1)
}
defun consumer(n) {
  var prod
  in(n, (? prod)) ; pattern
  consume-product(prod)
  consumer(n + 1)
}
```

Linda

Producers/Consumers is just as easy

```
defun producer(n) {
  out(n, make-product())
  producer(n + 1)
}
defun consumer(n) {
  var prod
  in(n, (? prod)) ; pattern
  consume-product(prod)
  consumer(n + 1)
}
```

We use a tag to ensure we consume values in the same order as they are produced (if that is important)

Linda

Exercise Think about the assessed coursework using Linda

Questions of granularity are just as important in Linda as elsewhere

Linda

Linda is easy to add to existing languages, usually as a library, occasionally with a minor tweak to the syntax for the patterns

Linda

Linda is easy to add to existing languages, usually as a library, occasionally with a minor tweak to the syntax for the patterns

Versions exist for C, Perl, Java and Prolog and others

Linda

Also

Linda

Also

- the blocking semantics lead to unwanted non-determinism:
see the `in` vying against the `read` above

Linda

Also

- the blocking semantics lead to unwanted non-determinism: see the `in` vying against the `read` above
- some implementations have non-blocking variants of `in` and `read`, but this just adds to the uncertainty

Linda

Also

- the blocking semantics lead to unwanted non-determinism: see the `in` vying against the `read` above
- some implementations have non-blocking variants of `in` and `read`, but this just adds to the uncertainty
- the low-level, unstructured nature of Linda can lead to awkward code: every application needs some mechanism to structure the tuples (tags being the simplest)

Linda

Also

- the blocking semantics lead to unwanted non-determinism: see the `in` vying against the `read` above
- some implementations have non-blocking variants of `in` and `read`, but this just adds to the uncertainty
- the low-level, unstructured nature of Linda can lead to awkward code: every application needs some mechanism to structure the tuples (tags being the simplest)
- there is no fairness on selecting tuples: a tuple can be ignored indefinitely if there are others that can be chosen

Linda

Also

- the blocking semantics lead to unwanted non-determinism: see the `in` vying against the `read` above
- some implementations have non-blocking variants of `in` and `read`, but this just adds to the uncertainty
- the low-level, unstructured nature of Linda can lead to awkward code: every application needs some mechanism to structure the tuples (tags being the simplest)
- there is no fairness on selecting tuples: a tuple can be ignored indefinitely if there are others that can be chosen
- junk can collect in the pool: tuples put in but never taken out. This can slow down the matching

Linda

Also

- the blocking semantics lead to unwanted non-determinism: see the `in` vying against the `read` above
- some implementations have non-blocking variants of `in` and `read`, but this just adds to the uncertainty
- the low-level, unstructured nature of Linda can lead to awkward code: every application needs some mechanism to structure the tuples (tags being the simplest)
- there is no fairness on selecting tuples: a tuple can be ignored indefinitely if there are others that can be chosen
- junk can collect in the pool: tuples put in but never taken out. This can slow down the matching
- the pool can be a bottleneck

Linda

Further

Linda

Further

- detecting when the program needs to terminate is a problem: this could be done by putting a special “end of program” tuple in the pool; but then threads have the overhead of constantly checking for that tuple (and you need a non-blocking read to do so). Or have an extra field in every tuple that is a status flag, etc.

Linda

Further

- detecting when the program needs to terminate is a problem: this could be done by putting a special “end of program” tuple in the pool; but then threads have the overhead of constantly checking for that tuple (and you need a non-blocking read to do so). Or have an extra field in every tuple that is a status flag, etc.
- *aliasing* is a problem: careful constructions of name schemes (tags, usually) are needed to ensure that tuples are not accidentally picked up by the wrong threads

Linda

Further

- detecting when the program needs to terminate is a problem: this could be done by putting a special “end of program” tuple in the pool; but then threads have the overhead of constantly checking for that tuple (and you need an non-blocking read to do so). Or have an extra field in every tuple that is a status flag, etc.
- *aliasing* is a problem: careful constructions of name schemes (tags, usually) are needed to ensure that tuples are not accidentally picked up by the wrong threads
- related is *temporal aliasing*, where information about the order tuples were put into the pool is lost: again an enumeration tag can fix this, but it has to be coded

Linda

So extensions of Linda exist, e.g., using multiple pools to structure, thus avoiding the first kind of aliasing

Linda

So extensions of Linda exist, e.g., using multiple pools to structure, thus avoiding the first kind of aliasing

Now pools become first-class objects, and you can pass pools via pools to other threads

Linda

So extensions of Linda exist, e.g., using multiple pools to structure, thus avoiding the first kind of aliasing

Now pools become first-class objects, and you can pass pools via pools to other threads

But, as always, this moves away from the initial simplicity of the Linda concept

Linda

Linda

Linda

Linda

- is a simple abstract model of parallelism

Linda

Linda

- is a simple abstract model of parallelism
- can map reasonably well to different kinds of hardware (shared and distributed)

Linda

Linda

- is a simple abstract model of parallelism
- can map reasonably well to different kinds of hardware (shared and distributed)
- is explicitly non-deterministic, with the non-determinism mostly well delineated

Linda

Linda

- is a simple abstract model of parallelism
- can map reasonably well to different kinds of hardware (shared and distributed)
- is explicitly non-deterministic, with the non-determinism mostly well delineated
- is not suited for all kinds of problem

Linda

Linda

- is a simple abstract model of parallelism
- can map reasonably well to different kinds of hardware (shared and distributed)
- is explicitly non-deterministic, with the non-determinism mostly well delineated
- is not suited for all kinds of problem
- is not widely used, but you do see Linda being mentioned now and again, mostly for coordination between other systems

Linda

For example, the computational chemistry (molecule simulation) package Gaussian uses OpenMP on a node and uses Linda between nodes