Unifying Math Ontologies: A Tale of Two Standards

Differentiating between analysis and algebra

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A thought

Whenever anyone says "you know what I mean", you can be pretty sure that *he* does not know what he means, for if he did, he would tell you.

— H. Davenport (1907–1969)

OpenMath and MathML share the goal of representing mathematics "as it is", rather than "as it ought to be". A relevant example of the difference is given by [Kamareddine & Nederpelt, 2004], where the original text is

The function
$$\sqrt{|x|}$$
 is not differentiable at 0 (1)

while its formalised equivalent is

$$\neg(\lambda_{x:\mathbf{R}}(\sqrt{|x|}) \text{ is differentiable at } 0).$$
 (2)

The key features are the typing of x as being in **R**, and the conversion of $\sqrt{|x|}$ from an expression to a function.

• what one learned in calculus/analysis about *functions*, which we will write as $D_{\epsilon\delta}$: the "differentiation of $\epsilon-\delta$ analysis" (similarly $\frac{d}{d_{\epsilon\delta}x}$, and its inverse $_{\epsilon\delta}\int$);

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- what is taught in differential algebra about (*expressions* in) differential fields, which *we* will write as D_{DA} : the "differentiation of differential algebra" (similarly $\frac{d}{d_{DA}x}$, and its inverse $_{DA} \int$).

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- (2) is unashamedly the former, while (1) talks about a function, but actually gives an expression.

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$$2x \neq 2y,\tag{3}$$

$$(\lambda x.2x) =$$
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studies four areas (which in fact turn out to be inter-related):constructions with bound variables;

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- (We shan't talk about the last in this presentation.)

MathML 2's rules on <bvar>

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Variables in bvar constructions 'bind' the corresponding variable occurrences in the scope of the parent of the bvar. However, the variable may (e.g. \forall) or may not (e.g. $\frac{d}{dx}$) be bound in the sense of α -convertibility.

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If there's a <condition>, its variables are as bound as the others.

$\lambda\text{-notation}$

λ -notation

To motivate the λ -notation, consider the everyday mathematical expression 'x - y'. This can be thought of as defining either a function f of x or g of y ... And there is need for a notation that gives f and g different names in some systematic way. In practice mathematicians usually avoid this need by various 'ad hoc' special notations, but these can get very clumsy when higher-order functions are involved. [Hindley & Seldin, 2008, p. 1]

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MathML 3 introduces a formal bind to take the guessing out of the MathML 2 'rule' quoted above.

Some uses of condition are OK: e.g. $\forall x \in \mathbf{R}p(x)$

<apply> <forall/> <bvar><ci>x</ci></bvar> <condition><apply><in/><ci>x</ci><reals/></apply></condition "p(x)"

</apply>

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```
<apply>
  <forall/>
  <bvar><ci>x</ci></bvar>
  <condition><apply><in/><ci>x</ci><reals/></apply></condi-
"p(x)"
</apply>
<OMBIND>
  <OMA>
    <OMS name="forallin" cd="quant3"/>
    <OMS name="R" cd="setname1"/>
  </OMA>
  <OMBVAR> <OMV name="x"/> </OMBVAR>
"p(x)"
</OMBIND>
```

Some uses of condition are not: e.g. $\forall x, y \in \mathbf{R} : x - y \neq 0.\frac{1}{x-y} \in \mathbf{R}$

```
<apply>
  <forall/>
  <bvar><ci>x</ci><ci>y</ci></bvar>
  <condition>
    <apply><and>
      <apply><in/><ci>x</ci><reals/></apply>
      <apply><in/><ci>y</ci><reals/></apply>
      "x\ne v"
      </apply>
  </condition>
  \frac{1}{x-y} \in \mathbb{R}
</apply>
```

<lowlimit> <cn>0</cn> </lowlimit> <uplimit> <ci>a</ci> </uplimit>

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<condition> <apply><in/> <ci>x</ci> <ci>D</ci> </apply> </condition>

<lowlimit> <cn>0</cn> </lowlimit> <uplimit> <ci>a</ci> </uplimit>

<condition> <apply><in/> <ci>x</ci> <ci>D</ci> </apply> </condition>

<domainofapplication>
 <ci>D</ci>
</domainofapplication>

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<condition> <apply><in/> <ci>x</ci> <ci>D</ci> </apply> </condition>

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<domainofapplication>
 <ci>D</ci>
</domainofapplication>

All equivalent in MathML 2. OpenMath can't easily model the second.

Is this a problem? consider multi-dimensional calculus

[Borwein & Erdelyi,1995, p. 189] has a real integral over a curve in the complex plane,

$$\frac{1}{2\pi} \int_{|t|=R} \left| \frac{f(t)}{t^{n+1}} \right| |dt| \tag{6}$$

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[Apostol, 1967, p. 413, exercise 4] has an integral where we clearly want to connect the variables in the integrand to the variables defining the set:

$$\int \int \int \int \left\{ \frac{x^2}{a^2} + \frac{y^2}{b^2} + \frac{z^2}{c^2} \le 1 \right\} \left\{ \frac{x^2}{a^2} + \frac{y^2}{b^2} + \frac{z^2}{c^2} \right\} dx dy dz$$
(7)

```
<OMBIND>
<OMA>
<OMS cd="calculus_new"
name="tripleintcond"/>
"\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}\le1"
</OMA>
<OMBVAR>"x,y,z"</OMBVAR>
"\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}"
</OMBIND>
```

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<OMS cd="calculus_new"
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<OMBVAR>"x,y,z"</OMBVAR>
"\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}"
</OMBIND>
```

Forbidden since the binder is not in its own scope.

```
<OMBIND>
<OMS cd="calculus_new"
name="tripleintcond"/>
<OMBVAR>"x,y,z"</OMBVAR>
"\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}\le1"
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"\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}"
</OMBIND>
```

Forbidden since the binder is only allowed one argument.

```
<OMBIND>
<OMS cd="calculus_new"
     name="tripleintcond"/>
<OMBVAR>"x,y,z"</OMBVAR>
<OMA>
   <OMS cd="calculus new"
       name="tripleint_inner"/>
   \frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}\
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   \frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}{c^2}
</OMA>
</OMBIND>
```

Legal, but unnatural.

Solution 4: bind separately

<OMA>

```
<OMS cd="calculus_new"
    name="tripleintcond"/>
    "\lambda{x,y,z}.\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}
    "\lambda{x,y,z}.\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}
</OMA>
```

<OMA> <OMS cd="calculus_new" name="tripleintco

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name="tripleintcond"/>
    "\lambda{x,y,z}.\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}
    "\lambda{x,y,z}.\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}
</OMA>
```

which is equivalent to

```
<OMA>
<OMS cd="calculus_new"
name="tripleintcond"/>
"\lambda{x,y,z}.\frac{x^2}{a^2}+\frac{y^2}{b^2}+\frac{z^2}
"\lambda{z,y,x}.\frac{z^2}{a^2}+\frac{y^2}{b^2}+\frac{x^2}
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Our proposal: legitimise 2

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Option 2 is our preferred route.

$$\forall x, y \in \mathbf{R} : x - y \neq 0. \frac{1}{x - y} \in \mathbf{R}$$
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<OMA>

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<OMS name="forallincond" cd="quant3"/>
<OMS name="R" cd="setname1">
</OMA>
<OMBVAR><OMV name="x"/><OMV name="y"/></OMBVAR>
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- This is most naturally done by extending OMBIND